

Yong Li





Shoaib Kamil





IVLA: Compilable Markdown for Linear Algebra

Alec Jacobson





University of Toronto



Yotam Gingold





BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI



Fig. 1. The skinned model (left) is produced directly from the "unrigged" rigid bind model using our Direct Delta Mush algorithm. DDM can produce equivalent results to the Delta Mush algorithm but uses a direct local computation rather than the iterated global "mush" runtime smoothing of DM. The DM and DDM algorithms both provide greatly simplified authoring. They do not have the bulge and cleft artifacts common to other methods, which are prominent in the under-arm and hip regions (respectively) in this example (red arrows). DDM offers further advantages over DM, as described in the paper.

A significant fraction of the world's population have experienced virtual characters through games and movies, and the possibility of online VR social experiences may greatly extend this audience. At present, the skin deformation for interactive and real-time characters is typically computed using geometric skinning methods. These methods are efficient and simple to implement, but obtaining quality results requires considerable manual "rigging" effort involving trial-and-error weight painting, the addition of virtual helper bones, etc. The recently introduced Delta Mush algorithm largely solves this rig authoring problem, but its iterative computational approach has prevented direct adoption in real-time engines.

This paper introduces Direct Delta Mush, a new algorithm that simultaneously improves on the efficiency and control of Delta Mush while generalizing previous algorithms. Specifically, we derive a direct rather than iterative algorithm that has the same ballpark computational form as some previous geometric weight blending algorithms. Straightforward variants of the algorithm are then proposed to further optimize computational and storage cost with insignificant quality losses. These variants are equivalent to special cases of several previous skinning algorithms.

Our algorithm simultaneously satisfies the goals of reasonable efficiency, quality, and ease of authoring. Further, its explicit decomposition of rotational and translational effects allows independent control over bending versus twisting deformation, as well as a skin sliding effect.

Authors' addresses: Binh Huy Le, SEED - Electronic Arts, Redwood City, CA, bbinh85@ gmail.com; JP Lewis, Google AI, San Francisco, CA, noisebrain@gmail.com.

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

CCS Concepts: • Computing methodologies → Animation.

Additional Key Words and Phrases: skinning, skeletal animation, delta mush, real time, deformation, character animation

ACM Reference Format:

Binh Huy Le and JP Lewis. 2019. Direct Delta Mush Skinning and Variants. *ACM Trans. Graph.* 38, 4, Article 113 (July 2019), 13 pages. https://doi.org/10. 1145/3306346.3322982

1 INTRODUCTION

Typically characters are the main focus of any movie or game. Major characters are often humans or animals, and thus are articulated models with rigid bones underlying deformable flesh and skin. Other objects in the scene such as trees can deform and may also be represented with a similar underlying approach. A key focus in all these cases is getting the deformation right.

A character deformation method suitable for games and interactive applications such as animation should have the following characteristics: (1) speed, (2) quality, (3) simplicity of setup and authoring. Existing approaches to character deformation can be very broadly classified into geometric skinning and simulation approaches. Simulation approaches produce the highest quality but may be less suitable in terms of criteria (1) and (3). Regarding speed, simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. T results t 1...n at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. ^{uivalent} d DDM $\frac{augoritin}{under-a}$ For convenience, we concatenate all vectors \mathbf{u}_i to a matrix $\mathbf{U} =$ $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}.$ A signif charact

The deformation of U is driven by a LBS model with m bones. ^{ta mush,} social er deforma using ge The transformation of bone j = 1..m is $M_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the to impl ^{*rigging} weight of bone *j* on vertex *i*. The weights are required to be affine, ^{Variants.} i. org/10. $\sum_{approac}^{largely}$ i.e. $\sum_{j=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect on our formulation. Note that we use the LBS model for the sake of eralizing generality, but rigid binding is more common in practice, i.e. each culated 1. Other ^{previou} vertex is only assigned to one bone ($w_{ij} \in \{0,1\}$). The skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as:

 $\mathbf{v}_i = \sum_{i=1}^m w_{ij} \mathbf{M}_j \mathbf{u}_i , \ i = 1..n$

d interlowing up and can be tion ap lity but g speed

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

Our a

quality,

tional

versus f

Authors

simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model 3.1

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. Tresults t 1..*n* at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. ^{uivalent} under-a For convenience, we concatenate all vectors \mathbf{u}_i to a matrix U = A signif $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}$. charact

The deformation of U is driven by a LBS model with *m* bones. ^{ta mush,} social er deforma using ge The transformation of bone j = 1..m is $M_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the to impl ^{*rigging} weight of bone j on vertex i. The weights are required to be affine, ^{Variants.} $\sum_{approac}^{largely}$ i.e. $\sum_{j=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect on our formulation. Note that we use the LBS model for the sake of eralizing generality, but rigid binding is more common in practice, i.e. each culated 1. Other ^{previou} vertex is only assigned to one bone ($w_{ij} \in \{0,1\}$). The skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as:

 $\mathbf{v}_i = \sum_{i=1}^n w_{ij} \mathbf{M}_j \mathbf{u}_i \,, \, i = 1..n$

d interlowing up and can be tion ap lity but g speed

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

Our a

quality,

tional

versus t

Authors

simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

10

11

12

13

14

15

16

```
#include <Eigen/Core>
   #include <Eigen/Dense>
   #include <Eigen/Sparse>
   #include <iostream>
   void get_skinned_geometry(const Eigen::MatrixXd & w,
                              const std::vector<Eigen::Matrix<double, 4, 4>> & M,
                              const Eigen::MatrixXd & U,
                              Eigen::MatrixXd & V)
       V.resize(4, U.cols());
       for (int i = 0; i < U.cols(); i++) {</pre>
            for (int j = 0; j < w.cols(); j++) {</pre>
                V.col(i) += w(i, j) * M[i] * U.col(i);
            }
17 }
```


Permission to make digital or hard copies of all or part of this work for personal or lassroom use is granted without fee provided that copies are not made or distribute for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model 3.1

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. Tresults t 1..*n* at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. ^{uivalent} under-a For convenience, we concatenate all vectors \mathbf{u}_i to a matrix U = A signif $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}$. charact

The deformation of U is driven by a LBS model with m bones. ^{ta mush,} social er deforma using ge The transformation of bone j = 1..m is $M_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the to impl ^{*rigging} weight of bone j on vertex i. The weights are required to be affine, ^{Variants.} $\sum_{approac}^{largely}$ i.e. $\sum_{j=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect on our formulation. Note that we use the LBS model for the sake of eralizing generality, but rigid binding is more common in practice, i.e. each culated 1. Other ^{previou} vertex is only assigned to one bone ($w_{ij} \in \{0,1\}$). The skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as:

 $\mathbf{v}_i = \sum_{i=1}^n w_{ij} \mathbf{M}_j \mathbf{u}_i \,, \, i = 1..n$

d interlowing up and can be tion ap lity but g speed

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

Our a

quality,

tional

versus t

Authors

simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

10

11

12

13

14

15

16

```
#include <Eigen/Core>
  #include <Eigen/Dense>
   #include <Eigen/Sparse>
   #include <iostream>
   void get_skinned_geometry(const Eigen::MatrixXd & w,
                              const std::vector<Eigen::Matrix<double, 4, 4>> & M,
                              const Eigen::MatrixXd & U,
                              Eigen::MatrixXd & V)
       V.resize(4, U.cols());
       for (int i = 0; i < U.cols(); i++) {</pre>
           for (int j = 0; j < w.cols(); j++) {</pre>
               V.col(i) += w(i, j) * M[i] * U.col(i);
           }
                            i should be j
17 }
```


Permission to make digital or hard copies of all or part of this work for personal or lassroom use is granted without fee provided that copies are not made or distribute for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model 3.1

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. Tresults t 1..n at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. In d DDM under-a For convenience, we concatenate all vectors \mathbf{u}_i to a matrix $\mathbf{U} = \mathbf{u}_i$ A signif $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}$.

The deformation of U is driven by a LBS model with *m* bones. ^{ta mush,} social er using ge The transformation of bone j = 1..m is $\mathbf{M}_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the to imple "rigging weight of bone j on vertex i. The weights are required to be affine, "variants. vi.org/10. $\sum_{approac}^{largely}$ i.e. $\sum_{j=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect This on our formulation. Note that we use the LBS model for the sake of eralizing generality, but rigid binding is more common in practice, i.e. each culated ı. Other previou vertex is only assigned to one bone $(w_{ij} \in \{0,1\})$. The skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as: Our a

 $\mathbf{v}_i = \sum_{i=1}^{n} w_{ij} \mathbf{M}_j \mathbf{u}_i , \ i = 1..n$

d interlowing up and can be tion aplity but g speed

Permission to make digital or hard copies of all or part of this work for personal or lassroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

quality,

versus t

Authors

simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

10

ACM Trans. Graph., Vol. 38, No. 4, Article 113. Publication date: July 2019.

盟

10

11

12

13

14

15

16

```
#include <Eigen/Core>
   #include <Eigen/Dense>
   #include <Eigen/Sparse>
   #include <iostream>
   void get_skinned_geometry(const Eigen::MatrixXd & w,
                              const std::vector<Eigen::Matrix<double, 4, 4>> & M,
                              const Eigen::MatrixXd & U,
                              Eigen::MatrixXd & V)
       V.resize(4, U.cols());
       for (int i = 0; i < U.cols(); i++) {</pre>
           for (int j = 0; j < w.cols(); j++) {</pre>
               V.col(i) += w(i, j) * M[i] * U.col(i);
                            i should be j
17 }
```

```
import numpy as np
def get_skinned_geometry(w, M, U):
     v = np.zeros((4, U.shape[1]))
     for i in range(0, U.shape[1]):
         for j in range(0, w.shape[1]):
             v[:, i] += w[i, j] * M[j] @ U[:, i]
     return v
Θ
```


BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model 3.1

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. results t 1..n at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. d DDM under-a For convenience, we concatenate all vectors \mathbf{u}_i to a matrix $\mathbf{U} = \mathbf{u}_i$ $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}.$ A signif

The deformation of U is driven by a LBS model with *m* bones. ¹ social er using ge The transformation of bone j = 1..m is $\mathbf{M}_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the to impl "rigging weight of bone j on vertex i. The weights are required to be affined, "variants. virtual 1 weight of bone j on vertex i. The weights are required to be affined at j i.org/10. i.e. $\sum_{i=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect largely approac on our formulation. Note that we use the LBS model for the set of . Major eralizing generality, but rigid binding is more common in practice e. each culated 1. Other previou vertex is only assigned to one bone $(w_{ij} \in \{0,1\})$ T skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as d inter-

Our a quality, versus f

Authors'

Permission to make digital or hard copies of all or part of this work for personal or lassroom use is granted without fee provided that copies are not made or distribute for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

 $\mathbf{v}_i = \sum_{i=1} w_{ij} \mathbf{M}_j \mathbf{u}_i , \ i = 1..n$ simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

10

11

12

17

18

盟

lowing

up and

lity but

9

10

1) tion ap

```
#include <Eigen/Core>
#include <Eigen/Dense>
#include <Eigen/Sparse>
#include <iostream>
void get_skinned_geometry(const Eigen::MatrixXd & w,
                       const std::vector<Eigen::Matrix<double, 4, 4>> & M,
                       const Eigen::MatrixXd & U,
                       Eigen::MatrixXd & V)
   V.resize(4, U.cols());
   for (int i = 0; i < U.cols(); i++) {</pre>
       for (int j = 0; j < w.cols(); j++) {</pre>
           V.col(i) += w(i, j) * M[i] * U.col(i);
                      i should be j
       port numpy as np
           t_skinned_geometry(w, M, U):
   ⇒def
              np.zeros((4, U.shape[1]))
                 in range(0, U.shape[1]):
                   j in range(0, w.shape[1]):
                    v[:, i] += w[i, j] * M[j] @ U[:, i]
          return v
   Θ
```


BINH HUY LE, SEED - Electronic Arts JP LEWIS, Google AI

Original Delta Mush Model 3.1

This section revises the original Delta Mush (DM) model [Mancewicz et al. 2014] and presents notations to setup DM on top of a Linear Blend Skinning (LBS) model [Magnenat-Thalmann et al. 1988].

Assume that our character model is represented by a polygonal mesh with *n* vertices. The homogeneous position of vertex i =Fig. 1. Tresults t 1..*n* at the rest pose is $\mathbf{u}_i \in \mathbb{R}^4$, where the 4th component is 1. ^{uivalent} under-a For convenience, we concatenate all vectors \mathbf{u}_i to a matrix $\mathbf{U} = \mathbf{u}_i$ $[\mathbf{u}_1, \mathbf{u}_2, ..., \mathbf{u}_n] \in \mathbb{R}^{4 \times n}.$ A signif

The deformation of U is driven by a LBS model with *m* bones. ¹⁴ social er using ge The transformation of bone j = 1..m is $\mathbf{M}_j \in \mathbb{R}^{4 \times 4}$. Let w_{ij} be the "rigging weight of bone j on vertex i. The weights are required to be affine , "variants. virtual 1 weight of bone j on vertex i. The weights are required to be affine , "variants." $\sum_{approac}^{largely}$ i.e. $\sum_{j=1}^{m} w_{ij} = 1, \forall i$. Non-negativity and sparseness have no effect on our formulation. Note that we use the LBS model for the set of . Major eralizing generality, but rigid binding is more common in practice e. each culated 1. Other previou vertex is only assigned to one bone $(w_{ij} \in \{0,1\})$ T skinned also be ıs in all storage geometry $\mathbf{V} = [\mathbf{v}_1, \mathbf{v}_2, ..., \mathbf{v}_n] \in \mathbb{R}^{4 \times n}$ is computed as d inter-

 $\mathbf{v}_i = \sum_{i=1}^{n} w_{ij} \mathbf{M}_j \mathbf{u}_i \,, \, i = 1..n$

Oura quality versus f

Authors

Permission to make digital or hard copies of all or part of this work for personal or oom use is granted without fee provided that copies are not made or distribute for profit or commercial advantage and that copies bear this notice and the full citation on the first page. Copyrights for components of this work owned by others than the author(s) must be honored. Abstracting with credit is permitted. To copy otherwise, or republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Request permissions from permissions@acm.org.

© 2019 Copyright held by the owner/author(s). Publication rights licensed to ACM. 0730-0301/2019/7-ART113 \$15.00

https://doi.org/10.1145/3306346.3322982

simulation effects are not justified when nearly the same effect can be produced with a cheaper method. It should be remembered that character deformation is just one of many things that must be computed within the frame interval at typical frame rates of 24fps (movie animation), 60fps (games) or 120fps (VR). Other tasks include various rendering steps, gameplay AI, collision detection, other types of physics, etc. Simulation approaches are also not ideal in terms of simplicity. The rig may require constructing additional components

10

11

12

17

18

#1

#i

毘

lowing

up and

lity but

9

10

1) tion ap

#include <Eigen/Core>

AN EIGENANALYSIS OF I_5

The Eigensystem of *I*₅ 4.1

We will now show that the eigensystem of any energy expressed solely in terms of I_5 can be written down in closed form. The I_5 invariant can be written in several forms,

$$I_5 = \|\mathbf{F}\mathbf{a}\|_2^2 = \mathbf{a}^T \mathbf{C}\mathbf{a} = \mathrm{tr}(\mathbf{C}\mathbf{A}),\tag{5}$$

where $\mathbf{A} = \mathbf{a}\mathbf{a}^T$ and $\|\cdot\|_2^2$ denotes the squared Euclidean norm. The PK1 and Hessian in 3D are

$$\frac{\partial I_5}{\partial \mathbf{F}} = 2\mathbf{F}\mathbf{A}$$
(6)
$$\frac{\partial^2 I_5}{\partial \mathbf{f}^2} = 2\begin{bmatrix} \mathbf{A}_{00}\mathbf{I}_{3\times3} & \mathbf{A}_{01}\mathbf{I}_{3\times3} & \mathbf{A}_{02}\mathbf{I}_{3\times3} \\ \mathbf{A}_{10}\mathbf{I}_{3\times3} & \mathbf{A}_{11}\mathbf{I}_{3\times3} & \mathbf{A}_{11}\mathbf{I}_{3\times3} \\ \mathbf{A}_{20}\mathbf{I}_{3\times3} & \mathbf{A}_{21}\mathbf{I}_{3\times3} & \mathbf{A}_{22}\mathbf{I}_{3\times3} \end{bmatrix} = 2\mathbf{H}_5,$$
(7)

where $I_{3\times 3}$ is a 3×3 identity matrix, and A_{ij} is the (i, j) scalar entry of A. (Appendix A shows the matrix explicitly.) Since Eqn. 7 is constant in a, it is straightforward to state its eigensystem in closed form. In 3D, it contains three identical non-zero eigenvalues, $\lambda_{0,1,2} = 2 \|\mathbf{a}\|_2^2$, and since fiber directions are usually normalized, this simplifies

ACM Trans. Graph., Vol. 38, No. 4, Article 69. Publication date: July 2019.

1	$A_{ij} = \{ 1 if (i,j) \in E \}$
2	0 otherwise
3	D_ii = ∑_j A_ij
4	$L = D^{-1}(D-A)$
5	
6	where
7	
8	$E \in \{\mathbb{Z} \times \mathbb{Z}\}$
9	A ∈ R^(n×n)
10	n E Z

C++

$A_{ij} = \{ 1 \text{ if } (i,j) \in E \}$ 0 otherwise 3 D_ii = ∑_j A_ij $L = D^{-1}(D-A)$ 5 6 where 7 8 $E \in \{\mathbb{Z} \times \mathbb{Z}\}$ 9 A ∈ ℝ^(n×n) 10 11 n∈ℤ 12 */ #include <Eigen/Core> 13 #include <Eigen/Dense> 14 #include <Eigen/Sparse> 15 #include <iostream> 16 17 #include <set> 18 struct myExpressionResultType { 19 20 Eigen::SparseMatrix<double> A; 21 Eigen::SparseMatrix<double> D; 22 Eigen::SparseMatrix<double> L; 23 myExpressionResultType(const Eig 24 const Eigen::SparseMo 25 const Eigen::SparseMo 26 : A(A), 27 D(D), 28 L(L) 29 {} 30 }; 31

Python

```
A_{ij} = \{ 1 \text{ if } (i,j) \in E \}
 2
 3
              0 otherwise
    D_ii = ∑_j A_ij
 4
    L = D^{-1}(D-A)
 5
 6
    where
 7
 8
    E \in \{\mathbb{Z} \times \mathbb{Z}\}
 9
    A ∈ ℝ^(n×n)
10
    n∈ℤ
11
12
     11 11 11
13
    import numpy as np
    import scipy
14
    import scipy.linalg
15
    from scipy import sparse
16
17
    from scipy.integrate import quad
18
    from scipy.optimize import minimize
19
20
21
    class myExpressionResultType:
22
         def __init__( self, A, D, L):
23
             self.A = A
24
             self.D = D
25
             self.L = L
26
27
    def myExpression(E, n):
28
29
         E = frozenset(E)
30
31
         assert all( len(el) == 2 for el
```

MATLAB

LaTeX

1	<pre>function output = myExpression(E, n)</pre>	1 \documentclass[12pt]{art
2	% output = myExpression(E, n)	<pre>2 \usepackage{mathdots}</pre>
3	%	3 \usepackage[bb=boondox]{
4	% $A_{ij} = \{ 1 \text{ if } (i,j) \in E \}$	4 \usepackage{mathtools}
5	% 0 otherwise	5 \usepackage{amssymb}
6	% D_ii = Σ_j A_ij	6 \usepackage{libertine}
7	% $L = D^{-1}(D-A)$	<pre>7 \DeclareMathOperator*{\a</pre>
8	%	<pre>8 \DeclareMathOperator*{\a</pre>
9	% where	9 \usepackage[paperheight=
10	%	10 \let\originalleft\left
11	% E ∈ {ℤ×ℤ}	11 \let\originalright\right
12	% A $\in \mathbb{R}^{(n \times n)}$	12 \renewcommand{\left}{\ma
13	% n∈ℤ	13 \renewcommand{\right}{\a
14	if nargin==0	14 - \begin{document}
15	<pre>warning('generating random i</pre>	15
16	[E, n] = generateRandomData	16 - \begin{center}
17	end	<pre>17 \resizebox{\textwidth}{!</pre>
18	<pre>function [E, n] = generateRandor</pre>	18 - {
19	n = randi(10);	<pre>19 - \begin{minipage}[c]{\tex</pre>
20	E = [];	20 - \begin{align*}
21	<pre>dim_2 = randi(10);</pre>	<pre>21 \mathit{A}_{\mathit{i},</pre>
22	for i = 1 :dim_2	<pre>22 \mathit{D}_{\mathit{i},\</pre>
23	<pre>E = [E;randi(10), randi(</pre>	23 $\mathbf{L} \& = \mathbf{L} $
24	end	24 \intertext{where}
25	end	<pre>25 \mathit{E} & \in \{\mathit</pre>
26		<pre>26 \mathit{A} & \in \mathbb</pre>
27	<pre>assert(size(E,2) == 2)</pre>	<pre>27 \mathit{n} & \in \mathbb</pre>
28	<pre>assert(numel(n) == 1);</pre>	28 \\
29		29 \end{align*}
30	$Aij_0 = zeros(2,0);$	30 \end{minipage}
31	$Avals_0 = zeros(1,0);$	31 }

Related work: Markup languages

- LaTeX [Goossens et al. 1994]
- Markdown [Gruber and Swartz 2004]
- AsciiMath [Jipsen 2005]
- MathML [W3C 2016]

MARK	DOWN
1 *	# Warning
2	The **gamma
	non-positive
	(n-1)!
3	
4	Derived by I
	part, the ga
5	
6	\$\$\Gamma(z)
	.\$\$
7	
8	The notation
	complex num
	integral co
	the second l
9	
10	
11	\$\$al
12	\Gamma(z·
13	&= \Bigl[-x
	dx //
14	&= \lim_{x\1
	e^{-0}\right
15	aligned

a function** is defined for all complex numbers except the ve integers. For any positive integer \$\$n\$\$, \$\$\Gamma(n) =

Daniel Bernoulli, for complex numbers with a positive real gamma function is defined via a convergent improper integral:

=
$$int_0^i x^{z-1} e^{-x}, dx, \ \quad \Re(z) > 0$$

on \$\$\Gamma (z)\$\$ is due to Legendre. If the real part of the mber \$\$z\$\$ is strictly positive (\$\$\Re (z)>0\$\$), then the onverges absolutely, and is known as the Euler integral of kind. Using integration by parts, one sees that:

```
\to \infty}\left(-x^z e^{-x}\right) - \left(-0^z
nt) + z\int_0^\infty x^{z-1} e^{-x}\, dx.
ed}$$
```

PREVIEW

Warning

The **gamma function** is defined for all complex numbers except the non-positive integers. For any positive integer $n, \Gamma(n) = (n - 1)!$.

Derived by Daniel Bernoulli, for complex numbers with a positive real part, the gamma function is defined via a convergent improper integral:

$$\Gamma(z) \;=\; \int_0^\infty \, x^{z-1} e^{-x} \, dx, \qquad \, \Re(z) \;>\; 0 \; .$$

The notation $\Gamma(z)$ is due to Legendre. If the real part of the complex number z is strictly positive ($\Re(z) > 0$), then the integral converges absolutely, and is known as the Euler integral of the second kind. Using integration by parts, one sees that:

$$egin{aligned} \Gamma(z\,+\,1) &=& \int_0^\infty x^z e^{-x}\,dx \ &=& \left[-x^z e^{-x}
ight]_0^\infty \,+\, \int_0^\infty \,z x^{z-1} e^{-x}\,dx \ &=& \lim_{x o\infty}\,\,\left(-x^z e^{-x}
ight)\,-\,\,\left(-0^z e^{-0}
ight)\,\,+\,z\,\int_0^\infty \,x^{z-1} e^{-x}\,dx. \end{aligned}$$

Related work: DSLs for graphics

- [Perlin 1985]
- [Hanrahan and Lawson 1990]
- SafeGI [Ou and Pellacini 2010]
- Halide [Ragan-Kelley et al. 2012]
- VizGen [Yang et al. 2016]
- Simit [Kjolstad et al. 2016]
- Ebb [Bernstein et al. 2016]
- Opt [Devito et al. 2017]
- Slang [He et al. 2018]
- [Preussner 2018]
- Taichi [Hu et al. 2019]
- [Geisler et al. 2020]
- Penrose [Ye et al. 2020]
- TEG [Bangaru et al. 2021]

Decoupling Algorithms from Schedules for Easy Optimization of Image Processing Pipelines

Abstract

Using existing programming tools, writing high-performance image processing code requires sacrificing readability, portability, and modularity. We argue that this is a consequence of conflating what computations define the *algorithm*, with decisions about *storage* and the order of computation. We refer to these latter two concerns as the schedule, including choices of tiling, fusion, recomputation vs. storage, vectorization, and parallelism.

We propose a representation for feed-forward imaging pipelines that separates the algorithm from its schedule, enabling highperformance without sacrificing code clarity. This decoupling simplifies the algorithm specification: images and intermediate buffers become functions over an infinite integer domain, with no explicit storage or boundary conditions. Imaging pipelines are compositions of functions. Programmers separately specify scheduling strategies for the various functions composing the algorithm, which allows them to efficiently explore different optimizations without changing the algorithmic code.

We demonstrate the power of this representation by expressing a range of recent image processing applications in an embedded domain specific language called Halide, and compiling them for ARM, x86, and GPUs. Our compiler targets SIMD units, multiple cores, and complex memory hierarchies. We demonstrate that it can handle algorithms such as a camera raw pipeline, the bilateral grid, fast local Laplacian filtering, and image segmentation. The algorithms expressed in our language are both shorter and faster than state-of-the-art implementations

CR Categories: I.3.6 [Computer Graphics]: Methodology and Techniques-Languages

Keywords: Image Processing, Compilers, Performance

1 Introduction

Computational photography algorithms require highly efficient implementations to be used in practice, especially on powerconstrained mobile devices. This is not a simple matter of programming in a low-level language like C. The performance difference between naive C and highly optimized C is often an order of magnitude. Unfortunately, optimization usually comes at the cost of programmer pain and code complexity, as computation must be reorganized to achieve memory efficiency and parallelism.

Jonathan Ragan-Kelley* Andrew Adams* Sylvain Paris[†] Marc Levoy[‡] Saman Amarasinghe* Frédo Durand* *MIT CSAIL [†]Adobe [‡]Stanford University - (a) Clean C++ : 9.94 ms per megapixel -

for (int y = 0; y < in.height(); y++)</pre> $\int_{0}^{1} \frac{1}{x} = 0; x < \text{in.width}(j; x+1) \\ \text{blurred}(x, y) = (\text{tmp}(x, y-1) + \text{tmp}(x, y) + \text{tmp}(x, y+1))/3;$

void blur(const Image &in, Image &blurred) {
 Image tmp(in.width(), in.height());

for (int y = 0; y < in.height(); y++)

(b) Fast C++ (for x86) : 0.90 ms per megapixel void fast_blur(const Image &in, Image &blurred) { n128i one_third = _mm_set1_epi16(21846);

for (int x = 0; x < in.width(); x++)
tmp(x, y) = (in(x-1, y) + in(x, y) + in(x+1, y))/3;</pre>

#pragma omp parallel for for (int yTile = 0; yTile < in.height(); yTile += 32) { _m128i a, b, c, sum, avg; _m128i tmp[(256/8)*(32+2)]; for (int xTile = 0; xTile < in.width(); xTile += 256) {</pre> $c = _mm_load_sil28((_ml28i*)(inPtr))$ sum = _mm_add_epi16(_mm_add_epi16(a, b), c); avg = _mm_mulhi_epi16(sum, one_third); _mm_store_si128(tmpPtr++, avg); inPtr += 8;

}}
fmpPtr = tmp;
for (int y = 0; y < 32; y++) {
 __m128i *outPtr = (__m128i *) (& (blurred(xTile, yTile+y)));
for (int x = 0; x < 256; x += 8) {
 __m128i *outPtr = (__0256) (0);
}</pre> a = _mm_load_si128(tmpPtr+(2*256)/8): b = _mm_load_sil28(tmpPtr+256/8)
c = _mm_load_sil28(tmpPtr++); sum = _mm_add_epi16(_mm_add_epi16(a, b), c) avg = _mm_mulhi_epi16(sum, one_third)
_mm_store_si128(outPtr++, avg);

- (c) Halide : 0.90 ms per megapixel

Func halide_blur(Func in) {

Func tmp, blurred; Var x, y, xi, yi; // The algorithm

tmp(x, y) = (in(x-1, y) + in(x, y) + in(x+1, y))/3;blurred(x, y) = (tmp(x, y-1) + tmp(x, y) + tmp(x, y+1))/3;// The schedule

blurred.tile(x, y, xi, yi, 256, 32) .vectorize(xi, 8).parallel(y); tmp.chunk(x).vectorize(x, 8);

return blurred;

Figure 1: The code at the top computes a 3×3 box filter using the composition of a 1×3 and a 3×1 box filter (a). Using vectorization. multithreading, tiling, and fusion, we can make this algorithm more than $10 \times$ faster on a quad-core x86 CPU (b). However, in doing so we've lost readability and portability. Our compiler separates the algorithm description from its schedule, achieving the same performance without making the same sacrifices (c). For the full details about how this test was carried out, see the supplemental material.

Taichi: A Language for High-Performance Computation on Spatially Sparse Data Structures

YUANMING HU, MIT CSAIL TZU-MAO LI, MIT CSAIL and UC Berkeley LUKE ANDERSON, MIT CSAIL JONATHAN RAGAN-KELLEY, UC Berkeley FRÉDO DURAND, MIT CSAIL

Fig. 1. (Top) We propose the Taichi programming language, which exposes a high-level interface for developing and processing spatially sparse multi-level data structures, and an optimizing compiler that automatically reduces data structure overhead. Programmers write code as if they are accessing dense voxels, while specifying the data arrangement independently. Our compiler automatically generates optimized, high-performance code tailored to the data structure. This results in concise code and better performance than highly-optimized reference implementations for various tasks. (Bottom) A fluid simulation using the material point method, where two liquid jets collide with each other, forming a thin sheet structure. We used a three-level sparse voxel grid with sizes 1³, 4³, 16³. Involved voxels are visualized in green. Both simulation and rendering are done using programs written in Taichi.

3D visual computing data are often spatially sparse. To exploit such sparsity, people have developed hierarchical sparse data structures, such as multilevel sparse voxel grids, particles, and 3D hash tables. However, developing and using these high-performance sparse data structures is challenging, due to their intrinsic complexity and overhead. We propose Taichi, a new data-oriented programming language for efficiently authoring, accessing, and maintaining such data structures. The language offers a high-level, data structure-agnostic interface for writing computation code. The user independently specifies the data structure. We provide several elementary components with different sparsity properties that can be arbitrarily composed to create a wide range of multi-level sparse data structures. This decoupling of data structures from computation makes it easy to experiment

Authors' addresses: Yuanming Hu, MIT CSAIL, yuanming@mit.edu; Tzu-Mao Li, MIT CSAIL and UC Berkeley, tzumao@berkeley.edu; Luke Anderson, MIT CSAIL, lukea@ mit.edu; Jonathan Ragan-Kelley, UC Berkeley, jrk@berkeley.edu; Frédo Durand, MIT CSAIL, fredo@mit.edu.

Permission to make digital or hard copies of part or all of this work for personal o classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation n the first page. Copyrights for third-party components of this work must be honored For all other uses contact the owner/author(s © 2019 Copyright held by the owner/author(s). 0730-0301/2019/11-ART201 https://doi.org/10.1145/3355089.3356506

with different data structures without changing computation code, and allows users to write computation as if they are working with a dense array. Our compiler then uses the semantics of the data structure and index analysis to automatically optimize for locality, remove redundant operations for coherent accesses, maintain sparsity and memory allocations, and generate efficient parallel and vectorized instructions for CPUs and GPUs. Our approach yields competitive performance on common computational kernels such as stencil applications, neighbor lookups, and particle scattering. We demonstrate our language by implementing simulation, rendering,

and vision tasks including a material point method simulation, finite element analysis, a multigrid Poisson solver for pressure projection, volumetric path tracing, and 3D convolution on sparse grids. Our computation-data structure decoupling allows us to quickly experiment with different data arrangements, and to develop high-performance data structures tailored for specific computational tasks. With $\frac{1}{10}$ th as many lines of code, we achieve 4.55× higher performance on average, compared to hand-optimized reference implementations.

 $\texttt{CCS}\ \texttt{Concepts:} \bullet \textbf{Software and its engineering} \to \textbf{Domain specific lan-}$ guages; \bullet Computing methodologies \rightarrow Parallel programming languages; Physical simulation.

Additional Key Words and Phrases: Sparse Data Structures, GPU Computing.

Halide [Ragan-Kelley et al. 2012]

Taichi [Hu et al. 2019]

Related work: Languages for numerical computing

- YALMIP [Löfberg 2004]
- Fortress [Allen et al. 2005]
- APL [Iverson 2007]
- BLAC [Spampinato and Püschel 2014]
- Julia [Bezanson et al. 2017]
- TACO [Kjolstad et al. 2017]
- GENO [Laue et al. 2019]

The Tensor Algebra Compiler 77:13 code-gen(index-expr, iv) # iv is the index variable let L = merge-lattice(index-expr, iv) 2 | for (int i = 0; i < B1_size, i++) { # initialize sparse pos variables int pB1 = (0 * B1_size) + i; int pC1 = (0 * C1_size) + i; int pA1 = (0 * A1_size) + i; for Dj in sparse-dimensions(L) emit "int pDj = Dj_pos[pDj-1];" 1 int pB2 = B2_pos[pB1] int pC2 = C2_pos[pC1] # while all merged dimensions have more values emit "while(until-any-exhausted(merged-dimensions(Lp))) { initialize sparse idx variables int iB = B2 idx[pB2] for Dj in sparse-dimensions(Lp) emit "int ivDj = Dj_idx[pDj];' int jC = C2_idx[pC2]; int j = min(jB, jC); int pA2 = (pA1 * A2_size) + j # merge sparse idx variables if (jB == j && jC == j) A[pA2] = B[pB2] + C[pC2]; else if (jB == j) "int iv = min(["ivDj," Dj in sparse-dimensions(Lp)]);" # compute dense pos variables A[pA2] = B[pB2] emit "int pDj = (pDj-1 * Dj_size) + iv;" else if (jC == j) A[pA2] = C[pC2] # compute expressions available at this loop level emit-available-expressions(index-expr, iv) # Section 6.2 if (jB == j) pB2++; if (jC == j) pC2++; # one case per lattice point below Lp 2 | while (pB2 < B2_pos[pB1+1]) { 3-4 | int j = B2_idx[pB2]; 5 | int pA2 = (pA1 * A2_size) + j; 7 | A[pA2] = B[pB2]; </pre> emit-reduction-compute() # Section 6.2 emit-index-assembly() # Section 6.3 pB2++; emit-compute() # Section 6.2 if result dimension Dj is accessed with iv emit "pDj++;" emit "}" while (pC2 < C2 pos[pC1+1]) { int j = C2_idx[pC2]; int pA2 = (pA1 * A2_size) + j; 5 | 7 | # conditionally increment the sparse pos variables A[pA2] = C[pC2]; pC2++; 8 | emit "if (ivDj == iv) pDj++; (a) Recursive algorithm to generate code for tensor expressions (b) Generated sparse matrix addition code

(c) Iteration graph for matrix addition (d) Dense merge lattice for *i* (e) Sparse merge lattice for *j*

Fig. 11. (a) Recursive code generation algorithm for tensor index notation. (b) Generated code for a 256×256 sparse matrix addition, $A_{ij} = B_{ij} + C_{ij}$, where B and C's formats are (dense_{d1}, sparse_{d2}) and A's format is $(\text{dense}_{d1}, \text{dense}_{d2})$. (c-e) Internal representations used to generate the code. The algorithm and generated code have matching labels. The generated code is simplified in four ways: the outer loop is a for loop, if statements are nested in else branches, and if and min statements are removed from the last two while loops.

The last free variable is special because its loop nest is where the code writes to the output tensor. Loops nested above it compute available expressions (emit-available-expressions), which help avoid redundant computations. Loops nested below it are reduction loops that add sub-computations to reduction variables (emit-reduction-compute). Finally, the last free variable's loop combines the temporaries prepared by other loops to compute the final expression (emit-compute).

Proc. ACM Program. Lang., Vol. 1, No. OOPSLA, Article 77. Publication date: October 2017.

TACO [Kjolstad et al. 2017]

Related work: Languages for proof-checking

- Agda [Norell 2007]
- Lean [de Moura et al. 2015]
- Cog [Team 2021]

- show a = b, from calc

theorem le.antisymm : \forall {a b : \mathbb{Z} }, a \leq b \rightarrow b \leq a \rightarrow a = b := take a b : \mathbb{Z} , assume (H_1 : a \leq b) (H_2 : b \leq a), obtain (n : N) (Hn : a + n = b), from le.elim H_1 , obtain (m : N) (Hm : b + m = a), from le.elim H_2 , have H_3 : a + of_nat (n + m) = a + 0, from ... -- suppressed rest of the proof due to space limitations have H₆ : n = 0, from nat.eq_zero_of_add_eq_zero_right H₅,

a = a + 0 : add_zero $... = a + n : H_6$... = b : Hn

Lean example

IVLA combines conventional syntax with unambiguous execution

CONVENTIONAL LINEAR ALGEBRA SYNTAX

<section-header><section-header><text><figure><figure><text><text><text></text></text></text></figure></figure></text></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><page-header><section-header><section-header><image/><image/><image/><image/><image/><text><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></page-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><page-header><text><text><image/><image/><image/></text></text></page-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><image/><image/><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>
<section-header><section-header><section-header><section-header><text><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><text><text><image/><image/><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><image/><image/><image/><image/><image/><image/><section-header><section-header><section-header><section-header><section-header><section-header><section-header><image/><text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<page-header><page-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></page-header></page-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><image/><image/><image/><image/><image/><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>
<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><page-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></page-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><image/><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><image/><image/><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><page-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></page-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><image/><image/><image/><image/><image/><image/><image/><image/><text><text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><text><text><image/><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><image/><image/><image/><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>
<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><section-header><section-header><section-header><section-header><text><text><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></text></text></section-header></section-header></section-header></section-header></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><page-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></page-header></page-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><image/><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><page-header><text><text><text><image/><image/><image/><image/><image/><image/><image/><image/><image/><text></text></text></text></text></page-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<page-header><page-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></page-header></page-header>
<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header>	<page-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></page-header>	<page-header><page-header><text><text><text><image/><image/><image/><image/><text></text></text></text></text></page-header></page-header>	<section-header><section-header><text><text><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></page-header></page-header>	<image/> <image/> <section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><image/><image/><text><text><text><text><text><text></text></text></text></text></text></text></page-header>	<page-header><page-header><page-header><text></text></page-header></page-header></page-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><image/><image/><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>
<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<image/> <image/> <image/> <section-header><section-header><section-header><image/><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><page-header><page-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></page-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><page-header><section-header><page-header><text><text><image/><image/><image/><image/><image/><image/><image/><image/><page-header><text><text></text></text></page-header></text></text></page-header></section-header></page-header></section-header></page-header></page-header>
<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<page-header><page-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><image/><image/><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><text><text><text><text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><text><text><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><section-header><page-header><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></page-header></section-header></section-header></section-header></section-header></page-header></page-header>	<section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header>
<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><image/><text><text><text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><text><image/><image/><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><text><text><image/><text><text><text><text><text><text></text></text></text></text></text></text></text></text></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><image/><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><image/><text><text><text><text><text><text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><image/><image/><text><text><text><text><text><text><text><text><text><text><text><text><text></text></text></text></text></text></text></text></text></text></text></text></text></text></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><section-header><section-header><text><image/><image/><image/><text></text></text></section-header></section-header></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><section-header><text><image/><image/><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header></section-header>	<section-header><section-header><section-header><text><text><text><text><text><text></text></text></text></text></text></text></section-header></section-header></section-header>	<page-header><page-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header><section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></section-header></page-header></page-header>	

<section-header><section-header><section-header><section-header><text><image><image><image><image><image><text><text><text><text><text><text>

<section-header><section-header><section-header><text><text><text><text><text><text><text><text><text><text>

We consider the second of t

$$\begin{aligned} L_{2} = L_{2} - L_{2} L_{2} = L_{2} + L_{$$

Analysis of all 1987 Equations at SIGGRAPH 2019

$$\mathcal{P}^* = \arg\min_{\{\mathcal{P}\}} \sum_{i=1}^{N} \mathcal{L}_{\text{TASK}}(f_{\text{ISP}}(\mathbf{I}_i; \mathcal{P}), \mathbf{T}_i), \qquad (1)$$

$$C_t = \frac{\sum_k w_{t,k} \alpha_{t,k} C_{t,k}}{\sum_k w_{t,k} \alpha_{t,k}}.$$
(8)

$$f_{\mathcal{M}}(\mathbf{x}) = \sum_{i} w_{i} f(\mathbf{x}|\Theta_{i})$$
(1)

$$\mathbf{D} = \begin{bmatrix} \mathbf{R}_{z}(-\psi) & 0\\ 0 & 1 \end{bmatrix} \begin{pmatrix} \begin{bmatrix} \hat{\mathbf{u}}\\ \hat{\psi} \end{bmatrix} - \begin{bmatrix} \mathbf{u}\\ \psi \end{bmatrix} \end{pmatrix}.$$
 (8)

$$\mathbf{s} = (\bar{\boldsymbol{\phi}}^T, \boldsymbol{\omega}^T, \mathbf{D}^T, \mathbf{I}^T)^T,$$

$$\overline{T}(\mathbf{x},\mathbf{y}) = e^{-\overline{\tau}(\mathbf{x},\mathbf{y})} = e^{-\int_0^y \overline{\mu}_t(\mathbf{x}-s\omega) \, \mathrm{d}s}$$

$$\vec{p}_i(t) = f(\vec{q}(t), \vec{r}_i, \ell)$$

$$\mathbf{L}(\boldsymbol{\theta}, \boldsymbol{\phi}) = \sum_{t \in \mathcal{T}} \left(\sum_{i \in C} \left\| \hat{R}_i^t - R_i^t \right\|_1 + \lambda \delta(z^t) \right),$$

(10)

$$\mathcal{P}^* = \arg\min_{\{\mathcal{P}\}} \sum_{i=1}^{N} \mathcal{L}_{\text{TASK}}(f_{\text{ISP}}(\mathbf{I}_i; \mathcal{P}), \mathbf{T}_i), \qquad (1)$$

$$C_{t} = \frac{\sum_{k} w_{t,k} \alpha_{t,k} C_{t,k}}{\sum_{k} w_{t,k} \alpha_{t,k}}.$$
(8)

$$f_{\mathcal{M}}(\mathbf{x}) = \sum_{i} \mathbf{w}_{i} f(\mathbf{x} | \Theta_{i})$$
(1)

$$\mathbf{D} = \begin{bmatrix} \mathbf{R}_{z}(-\psi) & 0\\ 0 & 1 \end{bmatrix} \begin{pmatrix} \begin{bmatrix} \hat{\mathbf{u}}\\ \hat{\psi} \end{bmatrix} - \begin{bmatrix} \mathbf{u}\\ \psi \end{bmatrix} \end{pmatrix}.$$
(8)

$$\mathbf{s} = (\bar{\boldsymbol{\phi}}^T, \boldsymbol{\omega}^T, \mathbf{D}^T, \mathbf{I}^T)^T,$$

$$\overline{T}(\mathbf{x},\mathbf{y}) = e^{-\overline{\tau}(\mathbf{x},\mathbf{y})} = e^{-\int_0^y \overline{\mu}_t(\mathbf{x}-\mathbf{s}\omega) \, \mathrm{d}s}$$

$$\vec{p}_i(t) = f(\vec{q}(t), \vec{r}_i, \ell)$$

$$\mathbf{L}(\boldsymbol{\theta}, \boldsymbol{\phi}) = \sum_{t \in \mathcal{T}} \left(\sum_{i \in C} \left\| \hat{R}_i^t - R_i^t \right\|_1 + \lambda \delta(z^t) \right),$$

(10)

$$\mathcal{P}^{*} = \arg\min_{\{\mathcal{P}\}} \sum_{i=1}^{N} \mathcal{L}_{\text{TASK}}(f_{\text{ISP}}(\mathbf{I}_{i};\mathcal{P}),\mathbf{T}_{i}), \qquad (1)$$

$$C_{t} = \frac{\sum_{k} w_{t,k} \alpha_{t,k} C_{t,k}}{\sum_{k} w_{t,k} \alpha_{t,k}}.$$
(8)

$$f_{\mathcal{M}}(\mathbf{x}) = \sum_{i} \mathbf{w}_{i} f(\mathbf{x} | \Theta_{i})$$
(1)

$$\mathbf{D} = \begin{bmatrix} \mathbf{R}_{z}(-\psi) & 0 \\ 0 & 1 \end{bmatrix} \begin{pmatrix} \hat{\mathbf{u}} \\ \hat{\psi} \end{pmatrix} - \begin{bmatrix} \mathbf{u} \\ \psi \end{bmatrix}$$
(8)

$$\mathbf{s} = (\boldsymbol{\phi}^T, \boldsymbol{\omega}^T, \mathbf{D}^T, \mathbf{I}^T)^T, \qquad ($$

$$\overline{T}(\mathbf{x},\mathbf{y}) = e^{-\overline{\tau}(\mathbf{x},\mathbf{y})} = e^{-\int_0^u \overline{\mu}_t (\mathbf{x} - \mathbf{s}\omega) \, \mathrm{d}s}$$

$$\vec{p}_i(t) = f(\vec{q}(t), \vec{r}_i, \ell)$$

$$\mathbf{L}(\theta, \phi) = \sum_{t \in \mathcal{T}} \left(\sum_{i \in C} \left\| \hat{R}_i^t - R_i^t \right\|_1 + \lambda \delta(z^t) \right),$$

(10)

 $\begin{aligned} \omega &= ABC \\ \hat{d} &= x^{\mathsf{T}} \omega^{\mathsf{T}} \omega x \end{aligned}$

where $A \in \mathbb{R}^{3\times n}$ $B \in \mathbb{R}^{n\times n}$ $C \in \mathbb{R}^{\infty}$

• Single-letter identifiers are encouraged

 $\omega = ABC$ $\hat{d} = x^{\mathsf{T}}\omega^{\mathsf{T}}\omega x$

where $A \in \mathbb{R}^{3\times n}$ $B \in \mathbb{R}^{(n\times m)}$ $C \in \mathbb{R}^{m\times 2}$ x ∈ ℝ²

- Single-letter identifiers are encouraged
- Juxtaposition is multiplication

 $\omega = ABC$ $\hat{d} = \mathbf{x}^{\mathsf{T}} \boldsymbol{\omega}^{\mathsf{T}} \boldsymbol{\omega} \mathbf{x}$

where A $\in \mathbb{R}^{3\times n}$ $B \in \mathbb{R}^{(n\times m)}$ $C \in \mathbb{R}^{m\times 2}$ **x** ∈ ℝ²

- Single-letter identifiers are encouraged
- Juxtaposition is multiplication
- Unicode

 $\omega = ABC$ $\hat{d} = \mathbf{x}^{\mathsf{T}} \boldsymbol{\omega}^{\mathsf{T}} \boldsymbol{\omega} \mathbf{x}$

where A $\in \mathbb{R}^{3\times n}$ $B \in \mathbb{R}^{(n\times m)}$ $C \in \mathbb{R}^{m\times 2}$ x ∈ ℝ²

- Single-letter identifiers are encouraged
- Juxtaposition is multiplication
- Unicode
- Variables cannot be re-defined

 $\omega = ABC$ $\hat{d} = x^{T}\omega^{T}\omega x$

where A $\in \mathbb{R}^{3xn}$ $B \in \mathbb{R}^{(n\times m)}$ $C \in \mathbb{R}^{m\times 2}$ $\mathbf{X} \in \mathbb{R}^2$

Variables in IVLA

- Single-letter identifiers are encouraged
- Juxtaposition is multiplication
- Unicode
- Variables cannot be re-defined
- Compatible matrix and vector dimensions are statically checked (compile-time, not run-time).

 $\omega = ABC$ $\hat{d} = \mathbf{x}^{\mathsf{T}} \boldsymbol{\omega}^{\mathsf{T}} \boldsymbol{\omega} \mathbf{x}$

where A $\in \mathbb{R}^{3\times n}$ $B \in \mathbb{R}^{(n\times m)}$ $C \in \mathbb{R}^{m\times 2}$ $\mathbf{X} \in \mathbb{R}^2$

Matrices in I ULA

Matrices in I LA

2D matrix definitions:

 $L = \begin{bmatrix} I & M+yx^T \end{bmatrix}$ MT 0] where $M \in \mathbb{R}^{(m \times n)}$ x ∈ ℝ^n y ∈ ℝ^m

Matrices in I LA

2D matrix definitions:









2D matrix definitions:



 $L = \begin{bmatrix} I & M+yx^T \end{bmatrix}$ MT 0] where $M \in \mathbb{R}^{(m \times n)}$ x ∈ ℝ^n y ∈ ℝ^m













2D matrix definitions:



 $L = \begin{bmatrix} I & M+yx^T \end{bmatrix}$ MT 0] where $M \in \mathbb{R}^{(m \times n)}$ x ∈ ℝ^n y ∈ ℝ^m

2D matrix definitions:

L = [2a 0 3 k+1] where $a \in \mathbb{R}$ k ∈ ℝ



 $L_{ij} = M_{ij} + 7y_{i}$ where $M \in \mathbb{R}^{(m \times n)}$ y ∈ ℝ^m



 $L_{ij} = \{ 1 \text{ if } (i,j) \in E \}$ **0** otherwise $L_{ii} = -\sum_{j \in I} for j != i) L_{i,j}$ where $E \in \{ \mathbb{Z} \times \mathbb{Z} \}$ $L \in \mathbb{R}^{(n \times n)}$ n ∈ ℤ

• 2D matrix definitions:

 $L = \begin{bmatrix} 2a & 0 \\ 3 & k+1 \end{bmatrix}$ where $a \in \mathbb{R}$ $k \in \mathbb{R}$

• Element-wise:

$$L_ij = M_ij + 7y$$

where
$$M \in \mathbb{R}^{m\times n}$$

$$y \in \mathbb{R}^m$$



L_ij = { 1 if (i,j) ∈ E 0 otherwise L_ii = -∑_(j for j != i) L_i,j where E ∈ { ℤ×ℤ } L ∈ ℝ^(n×n) n ∈ ℤ

i

2D matrix definitions:

L = [2a 0 3 k+1] where $a \in \mathbb{R}$ k ∈ ℝ



 $L_{ij} = M_{ij} + 7y_{i}$ where $M \in \mathbb{R}^{(m \times n)}$ y ∈ ℝ^m



 $L_{ij} = \{ 1 \text{ if } (i,j) \in E \}$ **0** otherwise $L_{ii} = -\sum_{j \in I} for j != i) L_{i,j}$ where $E \in \{ \mathbb{Z} \times \mathbb{Z} \}$ $L \in \mathbb{R}^{(n \times n)}$ n ∈ ℤ

2D matrix definitions:

L = [2a 0 3 k+1] where $a \in \mathbb{R}$ k ∈ ℝ



 $L_{ij} = M_{ij} + 7y_{i}$ where $M \in \mathbb{R}^{(m \times n)}$ $y \in \mathbb{R}^m$



 $L_{ij} = \{ 1 \text{ if } (i,j) \in E \}$ **0** otherwise $L_{ii} = -\sum_{j \in I} for j != i) L_{i,j}$ where $E \in \{ \mathbb{Z} \times \mathbb{Z} \}$ L ∈ ℝ^(n×n) n∈ℤ

• 2D matrix definitions:

 $L = \begin{bmatrix} 2a & 0 \\ 3 & k+1 \end{bmatrix}$ where $a \in \mathbb{R}$ $k \in \mathbb{R}$



L_ij = M_ij + 7y_i where $M \in \mathbb{R}^{(m \times n)}$ $y \in \mathbb{R}^m$





$$\mathbf{g} = \frac{\partial \mathbf{F}^{T}}{\partial \mathbf{u}} : \mathbf{P}(\mathbf{F}).$$
(1)

$$E(x, x^{t}, v^{t}) = \frac{1}{2}x^{T}Mx - x^{T}Mx^{p} + h^{2}W(x).$$

$$K_{ij} = \sum_{\widehat{C} \in \widehat{\mathcal{M}}} \int_{\mathbf{g}(\widehat{C})} \nabla \phi_i(\mathbf{x}) \cdot \nabla \phi_j(\mathbf{x}) \, \mathrm{d}\mathbf{x}, \qquad ($$

$$M_i = \mathcal{T}(I_i).$$

$$\mathcal{L}_{\text{relight}} = \mathbb{E}\left[w_2 \cdot \mathcal{P}(I^{\mathbb{R}}, I^{\star})\right].$$

$$E_{E'} = \sum_{i \in I'} \sum_{f \in F_i} \frac{A(f)}{A'} \max(0, n_f^i \cdot d_i)$$

1)

$$\varphi_{\mathcal{G}\to\mathcal{S}} = \max_{\mathbf{g}\in\mathcal{G}} \left[\min_{\mathbf{s}\in\mathcal{S}} \phi(\mathbf{s},\mathbf{g}) \right].$$

(3)
$$\tilde{\chi}(x) = \frac{1}{2} + \left(\frac{1}{2} - \frac{\alpha}{4}\right) \frac{d(x)}{w(x)} + \frac{\alpha d^3(x)}{4w^3(x)} + O(|w(x)|).$$



$$\mathbf{g} = \frac{\partial \mathbf{F}^{T}}{\partial \mathbf{u}} : \mathbf{P}(\mathbf{F})$$
(1)

$$E(x, x^t, v^t) = \frac{1}{2}x^T M x - x^T M x^p + h^2 W(x).$$

$$K_{ij} = \sum_{\widehat{C} \in \widehat{\mathcal{M}}} \int_{\mathbf{g}(\widehat{C})} \nabla \phi_i(\mathbf{x}) \cdot \nabla \phi_j(\mathbf{x}) \, \mathrm{d}\mathbf{x}, \qquad ($$

$$M_i = \mathcal{T}(I_i).$$

$$\mathcal{L}_{\text{relight}} = \mathbb{E}\left[w_2 \cdot \mathcal{P}(I^{\mathbb{R}}, I^{\star})\right].$$

$$E_{E'} = \sum_{i \in I'} \sum_{f \in F_i} \frac{A(f)}{A'} \max(0, n_f^i \cdot d_i)$$

1)

$$\varphi_{\mathcal{G}\to\mathcal{S}} = \max_{\mathbf{g}\in\mathcal{G}} \left[\min_{\mathbf{s}\in\mathcal{S}} \phi(\mathbf{s},\mathbf{g}) \right].$$

(3)
$$\tilde{\chi}(x) = \frac{1}{2} + \left(\frac{1}{2} - \frac{\alpha}{4}\right) \frac{d(x)}{w(x)} + \frac{\alpha d^3(x)}{4w^3(x)} + O(|w(x)|).$$



$$\mathbf{g} = \frac{\partial \mathbf{F}^{T}}{\partial \mathbf{u}} : \mathbf{P}(\mathbf{F})$$
(1)

$$E(x, x^t, v^t) = \frac{1}{2}x^T M x - x^T M x^p + h^2 W(x)$$

$$K_{ij} = \sum_{\widehat{C} \in \widehat{\mathcal{M}}} \int_{\mathbf{g}(\widehat{C})} \nabla \phi_i(\mathbf{x}) \cdot \nabla \phi_j(\mathbf{x}) \, \mathrm{d}\mathbf{x}, \qquad ($$

$$M_i = \mathcal{T}(I_i).$$

$$\mathcal{L}_{\text{relight}} = \mathbb{E}\left[w_2 \cdot \mathcal{P}(I^{\mathbb{R}}, I^{\star})\right].$$

$$E_{E'} = \sum_{i \in I'} \sum_{f \in F_i} \frac{A(f)}{A'} \max(0, n_f^i \cdot d_i)$$

(1)

$$\varphi_{\mathcal{G}\to\mathcal{S}} = \max_{\mathbf{g}\in\mathcal{G}} \left[\min_{\mathbf{s}\in\mathcal{S}} \phi(\mathbf{s},\mathbf{g}) \right].$$

(3)
$$\tilde{\chi}(x) = \frac{1}{2} + \left(\frac{1}{2} - \frac{\alpha}{4}\right) \frac{d(x)}{w(x)} + \frac{\alpha d^3(x)}{4w^3(x)} + O(|w(x)|).$$



$$\mathbf{g} = \frac{\partial \mathbf{F}^{T}}{\partial \mathbf{u}} : \mathbf{P}(\mathbf{F})$$
(1)

$$E(x, x^t, v^t) = \frac{1}{2}x^T M x - x^T M x^p + h^2 W(x).$$

$$K_{ij} = \sum_{\widehat{C} \in \widehat{\mathcal{M}}} \int_{g(\widehat{C})} \nabla \phi_i(\mathbf{x}) \cdot \nabla \phi_j(\mathbf{x}) \, \mathrm{d}\mathbf{x}, \qquad (\mathbf{x}) \in \widehat{\mathcal{M}} \, \mathrm{d}\mathbf{x}$$

$$M_i = \mathcal{T}(I_i)$$

$$\mathcal{L}_{\text{relight}} = \mathbb{E}\left[w_2 \cdot \mathcal{P}(I^R, I^{\star})\right].$$

$$E_{E'} = \sum_{i \in I'} \sum_{f \in F_i} \frac{A(f)}{A'} \max(0, n_f^i \cdot d_i)$$

1)

$$\varphi_{\mathcal{G}\to\mathcal{S}} = \max_{\mathbf{g}\in\mathcal{G}} \left| \min_{\mathbf{s}\in\mathcal{S}} \phi(\mathbf{s},\mathbf{g}) \right|.$$

(3)
$$\tilde{\chi}(x) = \frac{1}{2} + \left(\frac{1}{2} - \frac{\alpha}{4}\right) \frac{d(x)}{w(x)} + \frac{\alpha d^3(x)}{4w^3(x)} + O(|w(x)|).$$



$$p(\mathbf{x}, t) = c_i(\mathbf{x}) \cos \omega_i t + d_i(\mathbf{x}) \sin \omega_i t, \qquad (5)$$
$$= \sqrt{c_i^2(\mathbf{x}) + d_i^2(\mathbf{x})} \cos(\omega_i t + \varphi_i(\mathbf{x})) \qquad (6)$$

$$B_{i} = \begin{bmatrix} \cos(\omega_{i}T) & \sin(\omega_{i}T) \\ -\sin(\omega_{i}T) & \cos(\omega_{i}T) \end{bmatrix}.$$
 (15)

$$p_I(r,t) = \frac{k^2 \rho c D_S(t)}{4\pi r} \cos\theta e^{-ikr} \tag{1}$$

$$h(\varphi) = \frac{1}{\pi} \arctan\left(\Re \lambda_0 + 2\Re \sum_{l=1}^m \lambda_l \exp(-il\varphi)\right) + \frac{1}{2}.$$

$$E_M(\mathbf{R}, \mathbf{t}) = \left\| \arcsin\left(\frac{\mathbf{q}^\top \mathbf{t}_{\times} \mathbf{R} \mathbf{p}}{\|\mathbf{t}_{\times} \mathbf{R} \mathbf{p}\|}\right) \right\|_2$$

$$W_M D_I = W_N^M, \text{ where}$$
$$(W_N^M)_{i,j} = \frac{1}{2} \left(\frac{\mu^M(T_\alpha)}{\mu^N(T_\alpha)} \cot \alpha_{ij}^N + \frac{\mu^M(T_\beta)}{\mu^N(T_\beta)} \cot \beta_{ij}^N \right).$$

$$c_p(\alpha, \epsilon) = \frac{K_p e \sec^2(\alpha + \epsilon)}{d_p \sec^2(\phi_{\alpha, \epsilon})}$$
(10)

$$\alpha_{max} = \tan^{-1} \left(\frac{M w_m}{2d_f + D_e} \right), \tag{4}$$









$$p(\mathbf{x}, t) = c_i(\mathbf{x}) \cos \omega_i t + d_i(\mathbf{x}) \sin \omega_i t \qquad (5)$$
$$= \sqrt{c_i^2(\mathbf{x}) + d_i^2(\mathbf{x})} \cos(\omega_i t + \varphi_i(\mathbf{x})) \qquad (6)$$

$$B_{i} = \begin{bmatrix} \cos(\omega_{i}T) & \frac{\sin(\omega_{i}T)}{-\sin(\omega_{i}T)} & \frac{\sin(\omega_{i}T)}{\cos(\omega_{i}T)} \end{bmatrix}.$$
 (15)

$$p_I(r,t) = \frac{k^2 \rho c D_S(t)}{4\pi r} \cos\theta e^{-ikr} \tag{1}$$

$$h(\varphi) = \frac{1}{\pi} \arctan\left(\Re \lambda_0 + 2\Re \sum_{l=1}^m \lambda_l \exp(-il\varphi)\right) + \frac{1}{2}.$$

$$E_M(\mathbf{R}, \mathbf{t}) = \left\| \arcsin\left(\frac{\mathbf{q}^\top \mathbf{t}_{\times} \mathbf{R} \mathbf{p}}{\|\mathbf{t}_{\times} \mathbf{R} \mathbf{p}\|}\right) \right\|_2$$

$$W_M D_I = W_N^M, \text{ where}$$
$$(W_N^M)_{i,j} = \frac{1}{2} \left(\frac{\mu^M(T_\alpha)}{\mu^N(T_\alpha)} \cot \alpha_{ij}^N + \frac{\mu^M(T_\beta)}{\mu^N(T_\beta)} \cot \beta_{ij}^N \right).$$

$$c_p(\alpha, \epsilon) = \frac{K_p e \sec^2(\alpha + \epsilon)}{d_p \sec^2(\phi_{\alpha, \epsilon})}$$
(10)

$$\alpha_{max} = \tan^{-1} \left(\frac{M w_m}{2d_f + D_e} \right), \tag{4}$$









$$p(\mathbf{x}, t) = c_i(\mathbf{x}) \cos \omega_i t + d_i(\mathbf{x}) \sin \omega_i t$$

$$= \sqrt{c_i^2(\mathbf{x}) + d_i^2(\mathbf{x})} \cos(\omega_i t + \varphi_i(\mathbf{x}))$$
(6)

$$B_{i} = \begin{bmatrix} \cos(\omega_{i}T) & \sin(\omega_{i}T) \\ -\sin(\omega_{i}T) & \cos(\omega_{i}T) \end{bmatrix}.$$
 (15)

$$p_I(r,t) = \frac{k^2 \rho c D_S(t)}{4\pi r} cos\theta^{-ikr}$$
(1)

$$h(\varphi) = \frac{1}{\pi} \arctan\left(\Re \lambda_0 + 2\Re \sum_{l=1}^m \lambda_l \exp(-il\varphi)\right) + \frac{1}{2}.$$

$$E_M(\mathbf{R}, \mathbf{t}) = \left\| \arcsin\left(\frac{\mathbf{q}^\top \mathbf{t}_{\times} \mathbf{R} \mathbf{p}}{\|\mathbf{t}_{\times} \mathbf{R} \mathbf{p}\|}\right) \right\|_2$$

$$W_M D_I = W_N^M, \text{ where}$$
$$(W_N^M)_{i,j} = \frac{1}{2} \left(\frac{\mu^M(T_\alpha)}{\mu^N(T_\alpha)} \cot \alpha_{ij}^N + \frac{\mu^M(T_\beta)}{\mu^N(T_\beta)} \cot \beta_{ij}^N \right).$$

$$c_p(\alpha, \epsilon) = \frac{K_p e \sec^2(\alpha + \epsilon)}{d_p \sec^2(\phi_{\alpha, \epsilon})}$$
(10)

$$\alpha_{max} = \tan^{-1} \left(\frac{M w_m}{2d_f + D_e} \right), \tag{4}$$









$$p(\mathbf{x}, t) = c_i(\mathbf{x}) \cos \omega_i t + d_i(\mathbf{x}) \sin \omega_i t$$

$$= \sqrt{c_i^2(\mathbf{x}) + d_i^2(\mathbf{x})} \cos(\omega_i t + \varphi_i(\mathbf{x}))$$
(6)

$$B_{i} = \begin{bmatrix} \cos(\omega_{i}T) & \sin(\omega_{i}T) \\ -\sin(\omega_{i}T) & \cos(\omega_{i}T) \end{bmatrix}.$$
 (15)

$$p_I(r,t) = \frac{k^2 \rho c D_S(t)}{4\pi r} cos\theta^{-ikr}$$
(1)

$$h(\varphi) = \frac{1}{\pi} \arctan\left(\Re \lambda_0 + 2\Re \sum_{l=1}^m \lambda_l \exp(-il\varphi)\right) + \frac{1}{2}.$$

$$E_M(\mathbf{R}, \mathbf{t}) = \left\| \operatorname{arcsin} \left(\frac{\mathbf{q}^\top \mathbf{t}_{\times} \mathbf{R} \mathbf{p}}{\|\mathbf{t}_{\times} \mathbf{R} \mathbf{p}\|} \right) \right\|_2$$

$$W_M D_I = W_N^M, \text{ where}$$
$$(W_N^M)_{i,j} = \frac{1}{2} \left(\frac{\mu^M(T_\alpha)}{\mu^N(T_\alpha)} \cot \alpha_{ij}^N + \frac{\mu^M(T_\beta)}{\mu^N(T_\beta)} \cot \beta_{ij}^N \right).$$

$$c_p(\alpha, \epsilon) = \frac{K_p e^{\sec^2(\alpha + \epsilon)}}{d_p \sec^2(\phi_{\alpha, \epsilon})}$$
(10)

$$\alpha_{max} = \tan^{-1} \left(\frac{M w_m}{2d_f + D_e} \right), \tag{4}$$













a = f(2) + g()where $f \in \mathbb{R} \rightarrow \mathbb{R}^{3\times3}$ $g \in \emptyset \rightarrow \mathbb{R}^{3\times3}$



• Externally defined functions

- Built-in Functions
 - Directly used
 - Import from built-in libraries:
 trigonometric and linearalgebra

$$a = f(2) + g()$$

where
$$f \in \mathbb{R} \rightarrow \mathbb{R}^{3\times3}$$

$$g \in \emptyset \rightarrow \mathbb{R}^{3\times3}$$

from trigonometry: sin c = sin(a) + exp(b)where $a \in \mathbb{R}$ $b \in \mathbb{R}$



• Externally defined functions

- Built-in Functions
 - Directly used
 - Import from built-in libraries:
 trigonometric and linearalgebra

$$a = f(2) + g()$$

where
$$f \in \mathbb{R} \rightarrow \mathbb{R}^{3\times3}$$

$$g \in \emptyset \rightarrow \mathbb{R}^{3\times3}$$



$$k_t(x,y) := \frac{e^{-d(x,y)^2/4t}}{(4\pi t)^{n/2}} j(x,y)^{-1/2} \left(1 + \sum_{i=1}^{\infty} t^i \Phi_i(x,y)\right).$$

$$\varphi_{ij_a} := \sum_{p=0}^{a-1} \tilde{\theta}_i^{j_p, j_{p+1}}.$$
(8)

$$W_{p}(f_{0}, f_{1}) = \min_{\mathbf{P}} \sum_{i,j} c(x_{i}, y_{j}) P_{i,j}$$
(5)

$$e \cdot k_0 + \sum_{j=1}^4 k_j \cdot e^j, \qquad (1)$$

$$L_{s} = \sum_{t} \|q_{\lambda}^{t} - q_{\lambda}^{t-1}\|_{2}^{2},$$

$$L_{r} = \sum_{t} \sum_{k} \left\| \Pi q_{k}^{t} - u_{k}^{t} \right\|_{2}^{2} c_{k}^{t},$$

$$R = \sum_{i=1}^{k} ||h'_{i} - h_{i}||_{2}^{2} + \sum_{i=1}^{k} \sum_{j=i+1}^{k} ||h'_{i,j} - h_{i,j}||_{2}^{2},$$

$$\hat{\boldsymbol{\alpha}}_p = \sum_i w_{pi} \hat{\boldsymbol{\alpha}}_i. \tag{44}$$





$$k_t(x,y) := \frac{e^{-d(x,y)^2/4t}}{(4\pi t)^{n/2}} j(x,y)^{-1/2} \left(1 + \sum_{i=1}^{\infty} t^i \Phi_i(x,y)\right).$$

$$\varphi_{ij_a} := \sum_{p=0}^{a-1} \tilde{\theta}_i^{j_p, j_{p+1}}.$$
(8)

$$W_{p}(f_{0}, f_{1}) = \min_{\mathbf{P}} \sum_{i,j} c(x_{i}, y_{j}) P_{i,j}$$
 (5)

$$e \cdot k_0 + \sum_{j=1}^4 k_j \cdot e^j, \qquad (1)$$

$$L_s = \sum_t \left\| q_\lambda^t - q_\lambda^{t-1} \right\|_2^2,$$

$$L_{r} = \sum_{t} \sum_{k} \|\Pi q_{k}^{t} - u_{k}^{t}\|_{2}^{2} c_{k}^{t},$$

$$R = \sum_{i=1}^{k} ||h'_{i} - h_{i}||_{2}^{2} + \sum_{i=1}^{k} \sum_{j=i+1}^{k} ||h'_{i,j} - h_{i,j}||_{2}^{2},$$

$$\hat{\boldsymbol{\alpha}}_p = \sum_i w_{pi} \hat{\boldsymbol{\alpha}}_i. \tag{44}$$





$$k_t(x,y) := \frac{e^{-d(x,y)^2/4t}}{(4\pi t)^{n/2}} j(x,y)^{-1/2} \left(1 + \sum_{i=1}^{\infty} t^i \Phi_i(x,y)\right).$$

$$\varphi_{ij_a} := \sum_{p=0}^{a-1} \tilde{\theta}_i^{j_p, j_{p+1}}.$$
(8)

$$W_{p}(f_{0}, f_{1}) = \min_{\mathbf{P}} \sum_{i,j} c(x_{i}, y_{j}) P_{i,j}$$
 (5)

$$e \cdot k_0 + \sum_{j=1}^4 k_j \cdot e^j, \qquad (1)$$

$$L_s = \sum_t \left\| q_\lambda^t - q_\lambda^{t-1} \right\|_2^2,$$

$$L_{r} = \sum_{t} \sum_{k} \|\Pi q_{k}^{t} - u_{k}^{t}\|_{2}^{2} c_{k}^{t},$$

$$R = \sum_{i=1}^{k} ||h'_{i} - h_{i}||_{2}^{2} + \sum_{i=1}^{k} \sum_{j=i+1}^{k} ||h'_{i,j} - h_{i,j}||_{2}^{2},$$

$$\hat{\boldsymbol{\alpha}}_p = \sum_i w_{pi} \hat{\boldsymbol{\alpha}}_i. \tag{44}$$





$$k_t(x,y) := \frac{e^{-d(x,y)^2/4t}}{(4\pi t)^{n/2}} j(x,y)^{-1/2} \left(1 + \sum_{i=1}^{\infty} t^i \Phi_i(x,y)\right)$$

$$\varphi_{ij_a} := \sum_{p=0}^{a-1} \tilde{\theta}_i^{j_p, j_{p+1}}.$$
(8)

$$W_{p}(f_{0}, f_{1}) = \min_{\mathbf{P}} \sum_{i,j} c(x_{i}, y_{j}) P_{i,j}$$
 (5)

$$e \cdot k_0 + \sum_{j=1}^4 k_j \cdot e^j, \qquad (1)$$

$$L_s = \sum_t \left\| q_\lambda^t - q_\lambda^{t-1} \right\|_2^2,$$

$$L_{r} = \sum_{t} \sum_{k} \|\Pi q_{k}^{t} - u_{k}^{t}\|_{2}^{2} c_{k}^{t},$$

$$R = \sum_{i=1}^{k} ||h'_{i} - h_{i}||_{2}^{2} + \sum_{i=1}^{k} \sum_{j=i+1}^{k} ||h'_{i,j} - h_{i,j}||_{2}^{2},$$

$$\hat{\boldsymbol{\alpha}}_p = \sum_i w_{pi} \hat{\boldsymbol{\alpha}}_i. \tag{44}$$





 $\sum_{i} a_i b_i + c$

 $\sum_{i} a_{i}b_{i} + c \qquad (\sum_{i} a_{i}b_{i}) + c \\ (\sum_{i} a_{i}b_{i} + c)$

 $\sum_{i} a_{i}b_{i} + c \qquad (\sum_{i} a_{i}b_{i}) + c \\ (\sum_{i} a_{i}b_{i} + c)$

equations



Complex summation formulas from SIGGRAPH 2019

Complex summation formulas from SIGGRAPH 2019

$$L_{\text{total}} = \sum_{(I^S, I^L, I^{L'}) \in \mathcal{A}} L_{\text{rend}}(I^S, I^L) + L_{\text{adjust_rend}}(I^S, I^{L'}) + w_1 L_{\text{smooth}}(\mathbb{D})$$

$$(6)$$

$$\langle F \rangle^{CV} = \langle F \rangle + \sum_{i=1}^{K} \gamma_i (G_i - \langle G_i \rangle)$$

$$= \sum_{i=1}^{K} \gamma_i G_i + \langle F \rangle - \sum_{i=1}^{K} \gamma_i \langle G_i \rangle$$

$$(14)$$

$$C_{\mathbf{v}_{1},\mathbf{v}_{2}}^{\mathbf{i}_{1},\mathbf{i}_{2}} \approx \frac{1}{N} \sum_{n=1}^{N} u_{\mathbf{v}_{1}}^{\mathbf{i}_{1},O^{n}} \cdot u_{\mathbf{v}_{2}}^{\mathbf{i}_{2},O^{n*}} - m_{\mathbf{v}_{1}}^{\mathbf{i}_{1}} \cdot m_{\mathbf{v}_{2}}^{\mathbf{i}_{2}*}.$$
 (6)

$$\tilde{u}(\mathbf{x},k) = \sum_{j} a_{jk} F_k(\mathbf{x} - \mathbf{y}_j, k) + u_{in}(\mathbf{x},k)$$
(22)
$$= -\sum_{j} a_{jk} \frac{i}{4} H_0^{(2)}(k||\mathbf{x} - \mathbf{y}_j||) + u_{in}(\mathbf{x},k).$$
(23)

 $\min_{\mathbf{a}} \qquad \|\ddot{\mathbf{q}}_{\mathrm{d}}(\mathbf{u}) - \ddot{\mathbf{q}}(\mathbf{a})\|^{2} + w_{\mathrm{reg}} \|\mathbf{a}\|^{2}$ subject to $\mathbf{M}\ddot{\mathbf{q}} + \mathbf{c} = \sum_{m} \mathbf{J}_{m}^{\top} \mathbf{f}_{m}(a_{m}) + \mathbf{J}_{\mathrm{c}}^{\top} \mathbf{f}_{\mathrm{c}} + \tau_{\mathrm{ext}} \qquad (13)$ $0 \le a_{m} \le 1 \quad \text{for} \quad \forall m.$

$$\mathcal{E}_{symm} \approx \sum_{i} \left[\cos \theta \left(p_{i} - q_{i} \right) \cdot n_{i} + \frac{1}{2} \cos \theta \left(\tilde{a} \times \left(p_{i} + q_{i} \right) \right) \cdot n_{i} + t \cdot n_{i} \right]^{2} \\ = \sum_{i} \cos^{2} \theta \left[\left(p_{i} - q_{i} \right) \cdot n_{i} + \left(\left(p_{i} + q_{i} \right) \times n_{i} \right) \cdot \tilde{a} + n_{i} \cdot \tilde{t} \right]^{2}, \qquad (9)$$

$$\mathcal{E}_{two-plane} = \sum_{i} \left[(Rp_{i} - R^{-1}q_{i} + t) \cdot (Rn_{p,i}) \right]^{2} + (14) \left[(Rp_{i} - R^{-1}q_{i} + t) \cdot (R^{-1}n_{q,i}) \right]^{2} \right]^{2}$$

$$f_{\mathbf{s},\mathbf{g}}(\mathbf{x}) = \sum_{i} a_{i} \phi(\mathbf{x}, \mathbf{x}_{i}) + \sum_{i} \mathbf{b}_{i}^{T} D^{0,1} \phi(\mathbf{x}, \mathbf{x}_{i}) + \mathbf{c}^{T} \mathbf{x} + d \qquad (3)$$

$$L(\hat{x}, x_{\Gamma}) = \sum_{j \in [1, s]} \left(E_j(x_j) + \frac{1}{2} \| z_j - R_{\Gamma_j} x_{\Gamma} \|_{K_j}^2 \right).$$
(5)

$$\Sigma_{J} = \left[\breve{q}_{J}\right]^{-1} = \left[\sum_{i} \alpha_{i} \breve{q}_{i}\right]^{-1} = \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1}\right)^{-1},$$

$$\mu_{J} = \Sigma_{J} \left(\sum_{i} \alpha_{i} \bar{q}_{i}\right) = \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1}\right)^{-1} \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1} \mu_{i}\right),$$
(13)
Complex summation formulas from SIGGRAPH 2019

$$L_{\text{total}} = \sum_{(I^{S}, I^{L}, I^{L'}) \in \mathcal{A}} L_{\text{rend}}(I^{S}, I^{L}) + L_{\text{adjust_rend}}(I^{S}, I^{L'}) + w_{1}L_{\text{smooth}}(\mathbb{D})$$

$$(6)$$

$$\langle F \rangle^{CV} = \langle F \rangle + \sum_{i=1}^{K} \gamma_i (G_i - \langle G_i \rangle)$$

$$= \sum_{i=1}^{K} \gamma_i G_i + \langle F \rangle - \sum_{i=1}^{K} \gamma_i \langle G_i \rangle$$

$$(14)$$

$$C_{\mathbf{v}_{1},\mathbf{v}_{2}}^{\mathbf{i}_{1},\mathbf{i}_{2}} \approx \frac{1}{N} \sum_{n=1}^{N} u_{\mathbf{v}_{1}}^{\mathbf{i}_{1},O^{n}} \cdot u_{\mathbf{v}_{2}}^{\mathbf{i}_{2},O^{n*}} - m_{\mathbf{v}_{1}}^{\mathbf{i}_{1}} \cdot m_{\mathbf{v}_{2}}^{\mathbf{i}_{2}*}.$$
 (6)

$$\tilde{u}(\mathbf{x},k) = \sum_{j} a_{jk} F_k(\mathbf{x} - \mathbf{y}_j, k) + u_{in}(\mathbf{x},k)$$
(22)
$$= -\sum_{j} a_{jk} \frac{i}{4} H_0^{(2)}(k||\mathbf{x} - \mathbf{y}_j||) + u_{in}(\mathbf{x},k).$$
(23)

 $\min_{\mathbf{a}} \qquad \|\ddot{\mathbf{q}}_{\mathrm{d}}(\mathbf{u}) - \ddot{\mathbf{q}}(\mathbf{a})\|^{2} + w_{\mathrm{reg}} \|\mathbf{a}\|^{2}$ $\mathrm{subject to} \qquad \mathbf{M}\ddot{\mathbf{q}} + \mathbf{c} = \sum_{m} \mathbf{J}_{m}^{\top} \mathbf{f}_{m}(a_{m}) + \mathbf{J}_{\mathrm{c}}^{\top} \mathbf{f}_{\mathrm{c}} + \tau_{\mathrm{ext}} \qquad (13)$ $0 \le a_{m} \le 1 \quad \text{for} \quad \forall m.$

$$\mathcal{E}_{symm} \approx \sum_{i} \left[\cos \theta \left(p_{i} - q_{i} \right) \cdot n_{i} + \frac{1}{2} \cos \theta \left(\tilde{a} \times \left(p_{i} + q_{i} \right) \right) \cdot n_{i} + t \cdot n_{i} \right]^{2} \\ = \sum_{i} \cos^{2} \theta \left[\left(p_{i} - q_{i} \right) \cdot n_{i} + \left(\left(p_{i} + q_{i} \right) \times n_{i} \right) \cdot \tilde{a} + n_{i} \cdot \tilde{t} \right]^{2}, \qquad (9)$$

$$\mathcal{E}_{two-plane} = \sum_{i} \left[(Rp_{i} - R^{-1}q_{i} + t) \cdot (Rn_{p,i}) \right]^{2} + (14) \left[(Rp_{i} - R^{-1}q_{i} + t) \cdot (R^{-1}n_{q,i}) \right]^{2} \right]^{2}$$

$$f_{\mathbf{s},\mathbf{g}}(\mathbf{x}) = \sum_{i} a_{i} \phi(\mathbf{x}, \mathbf{x}_{i}) + \sum_{i} \mathbf{b}_{i}^{T} D^{0,1} \phi(\mathbf{x}, \mathbf{x}_{i}) + \mathbf{c}^{T} \mathbf{x} + d \qquad (3)$$

$$L(\hat{x}, x_{\Gamma}) = \sum_{j \in [1, s]} \left(E_j(x_j) + \frac{1}{2} \| z_j - R_{\Gamma_j} x_{\Gamma} \|_{K_j}^2 \right).$$
(5)

$$\Sigma_{J} = \left[\breve{q}_{J}\right]^{-1} = \left[\sum_{i} \alpha_{i} \breve{q}_{i}\right]^{-1} = \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1}\right)^{-1},$$

$$\mu_{J} = \Sigma_{J} \left(\sum_{i} \alpha_{i} \bar{q}_{i}\right) = \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1}\right)^{-1} \left(\sum_{i} \alpha_{i} \Sigma_{i}^{-1} \mu_{i}\right),$$
(13)

Conservative rather than greedy



Conservative rather than greedy

Bounds inferred from index use

$$\sum_{i} a_{i}b_{i} + c \qquad (\sum_{i} a_{i}b_{i}) + c \\ (\sum_{i} a_{i}b_{i} + c)$$

$$\sum_{i \in \mathbb{N}} T_i p$$
where
$$w_i \in \mathbb{R}$$

$$T_i \in \mathbb{R}^{(4\times 4)}$$

$$p \in \mathbb{R}^{4}$$

Conservative rather than greedy

Bounds inferred from index use

$$\sum_{i} a_{i}b_{i} + c \qquad (\sum_{i} a_{i}b_{i}) + c \\ (\sum_{i} a_{i}b_{i} + c)$$

$$\sum_{i \in W_i} T_i p$$
where
$$w_i \in \mathbb{R}$$

$$T_i \in \mathbb{R}^{(4\times 4)}$$

$$p \in \mathbb{R}^4$$

Conservative rather than greedy

Bounds inferred from index use

$$\sum_{i} a_{i}b_{i} + c \qquad (\sum_{i} a_{i}b_{i}) + c \\ (\sum_{i} a_{i}b_{i} + c)$$

$$\sum_{i \in W_i \in T_i p}$$
where
$$w_i \in \mathbb{R}$$

$$T_i \in \mathbb{R}^{(4\times 4)}$$

$$p \in \mathbb{R}^4$$

$$\iint_{\Omega} \left(\frac{\rho}{\Delta t} \| \mathbf{u} - \mathbf{u}^* \|_2^2 + \mu \| \nabla \mathbf{u} \|_F^2 \right) dV, \tag{5}$$

$$r_k = ||\mathbf{A}\mathbf{x} - \mathbf{b}||_{\mathcal{K}_k}$$

$$\min_{\alpha} \| \sum_{i} \alpha_{i} E^{X_{i}} - E^{V} \|_{F}^{2} + \tau \| \alpha \|_{1}, \qquad (10)$$

(69)

$$\min_{\alpha} \| E^U C - C \sum_{i} \alpha_i E^{V_i} \|^2 + \tau \| \alpha \|_1.$$
(14)

$$E(\widetilde{\mathsf{L}}) = \|\mathsf{P}\mathsf{M}^{-1}\mathsf{L}\mathsf{I} - \widetilde{\mathsf{M}}^{-1}\widetilde{\mathsf{L}}\mathsf{P}\mathsf{I}\|_{\widetilde{\mathsf{M}}}^{2},$$
(3)

$$\min_{\{\mathbf{a}_l,\mathbf{b}_l,\mathbf{c}_l\}} \|R - \sum_{l=1}^{\overline{m}} \mathbf{a}_l \otimes \mathbf{b}_l \otimes \mathbf{c}_l\|_{\mathcal{F}}^2$$

$$\gamma \sum_{i=1}^{r} A_1(p_i) \|X_{12}(p_i,:) - X_2(q_i)\|_{M_2}^2 + A_2(q_i) \|X_{21}(q_i,:) - X_1(p_i)\|_{M_2}^2$$

$$\mathbf{w}_{j} = \left(\mathbf{R}_{j}^{W}\right)^{\top} \left(\frac{\mathbf{P} - \mathbf{t}_{j}^{W}}{\|\mathbf{P} - \mathbf{t}_{j}^{W}\|_{2}}\right).$$
(13)





$$\iiint_{\Omega} \left(\frac{\rho}{\Delta t} \| \mathbf{u} - \mathbf{u}^* \|_2^2 + \mu \| \nabla \mathbf{u} \|_F^2 \right) dV,$$
 (5)

$$r_k = ||\mathbf{A}\mathbf{x} - \mathbf{b}||_{\mathcal{K}_k} \tag{69}$$

$$\min_{\alpha} \| \sum_{i} \alpha_{i} E^{X_{i}} - E^{V} \|_{F}^{2} + \tau \| \alpha \|_{1}, \qquad (10)$$

$$\min_{\alpha} \| E^U C - C \sum_{i} \alpha_i E^{V_i} \|^2 + \tau \| \alpha \|_1.$$
(14)

$$E(\widetilde{\mathsf{L}}) = \|\mathsf{P}\mathsf{M}^{-1}\mathsf{L}\mathsf{I} - \widetilde{\mathsf{M}}^{-1}\widetilde{\mathsf{L}}\mathsf{P}\mathsf{I}\|_{\widetilde{\mathsf{M}}}^{2},$$
(3)

$$\min_{\{\mathbf{a}_l,\mathbf{b}_l,\mathbf{c}_l\}} \|R - \sum_{l=1}^{\overline{m}} \mathbf{a}_l \otimes \mathbf{b}_l \otimes \mathbf{c}_l\|_{\mathcal{F}}^2$$

$$\gamma \sum_{i=1}^{r} A_1(p_i) \|X_{12}(p_i,:) - X_2(q_i)\|_{M_2}^2 + A_2(q_i) \|X_{21}(q_i,:) - X_1(p_i)\|_{M_2}^2$$

$$\mathbf{w}_{j} = \left(\mathbf{R}_{j}^{W}\right)^{\top} \begin{pmatrix} \mathbf{P} - \mathbf{t}_{j}^{W} \\ \|\mathbf{P} - \mathbf{t}_{j}^{W}\|_{2} \end{pmatrix}.$$
 (13)





$$\iiint_{\Omega} \left(\frac{\rho}{\Delta t} \| \mathbf{u} - \mathbf{u}^* \|_2^2 + \mu \| \nabla \mathbf{u} \|_F^2 \right) dV,$$
 (5)

$$r_k = ||\mathbf{A}\mathbf{x} - \mathbf{b}||_{\mathcal{K}_k} \tag{69}$$

$$\min_{\alpha} \| \sum_{i} \alpha_{i} E^{X_{i}} - E^{V} \|_{F}^{2} + \tau \| \alpha \|_{1}$$
(10)

$$\min_{\alpha} \| E^{U}C - C \sum_{i} \alpha_{i} E^{V_{i}} \|^{2} + \tau \| \alpha \|_{1}$$
(14)

$$E(\widetilde{\mathsf{L}}) = \|\mathsf{P}\mathsf{M}^{-1}\mathsf{L}\mathsf{I} - \widetilde{\mathsf{M}}^{-1}\widetilde{\mathsf{L}}\mathsf{P}\mathsf{I}\|_{\widetilde{\mathsf{M}}}^{2},$$
(3)

$$\min_{\{\mathbf{a}_l,\mathbf{b}_l,\mathbf{c}_l\}} \|R - \sum_{l=1}^{\overline{m}} \mathbf{a}_l \otimes \mathbf{b}_l \otimes \mathbf{c}_l\|_{\mathcal{F}}^2$$

$$\gamma \sum_{i=1}^{r} A_1(p_i) \|X_{12}(p_i,:) - X_2(q_i)\|_{M_2}^2 + A_2(q_i) \|X_{21}(q_i,:) - X_1(p_i)\|_{M_2}^2$$

$$\mathbf{w}_{j} = \left(\mathbf{R}_{j}^{W}\right)^{\top} \begin{pmatrix} \mathbf{P} - \mathbf{t}_{j}^{W} \\ \|\mathbf{P} - \mathbf{t}_{j}^{W}\|_{2} \end{pmatrix}.$$
 (13)





$$\iint_{\Omega} \left(\frac{\rho}{\Delta t} \| \mathbf{u} - \mathbf{u}^* \|_2^2 + \mu \| \nabla \mathbf{u} \|_F^2 \right) dV, \tag{5}$$

$$r_k = ||\mathbf{A}\mathbf{x} - \mathbf{b}||_{\mathcal{K}_k} \tag{69}$$

$$\min_{\alpha} \|\sum_{i} \alpha_{i} E^{X_{i}} - E^{V} \|_{F}^{2} + \tau \|\alpha\|_{1}$$
(10)

$$\min_{\alpha} \|E^{U}C - C\sum_{i} \alpha_{i} E^{V_{i}}\|^{2} + \tau \|\alpha\|_{1}$$
(14)

$$E(\widetilde{\mathsf{L}}) = \|\mathsf{P}\mathsf{M}^{-1}\mathsf{L}\mathsf{I} - \widetilde{\mathsf{M}}^{-1}\widetilde{\mathsf{L}}\mathsf{P}\mathsf{I}\|_{\widetilde{\mathsf{M}}}^{2},$$
(3)

$$\min_{\{\mathbf{a}_l,\mathbf{b}_l,\mathbf{c}_l\}} \|R - \sum_{l=1}^m \mathbf{a}_l \otimes \mathbf{b}_l \otimes \mathbf{c}_l\|_{\mathcal{F}}^2$$

$$\gamma \sum_{i=1}^{r} A_1(p_i) \|X_{12}(p_i,:) - X_2(q_i)\|_{M_2}^2 + A_2(q_i) \|X_{21}(q_i,:) - X_1(p_i)\|_{M_2}^2$$

$$\mathbf{w}_{j} = \left(\mathbf{R}_{j}^{W}\right)^{\top} \begin{pmatrix} \mathbf{P} - \mathbf{t}_{j}^{W} \\ \|\mathbf{P} - \mathbf{t}_{j}^{W}\|_{2} \end{pmatrix}.$$
 (13)





$$\iint_{\Omega} \left(\frac{\rho}{\Delta t} \| \mathbf{u} - \mathbf{u}^* \|_2^2 + \mu \| \nabla \mathbf{u} \|_F^2 \right) dV, \tag{5}$$

$$r_k = ||\mathbf{A}\mathbf{x} - \mathbf{b}||_{\mathcal{K}_k}$$
(69)

$$\min_{\alpha} \| \sum_{i} \alpha_{i} E^{X_{i}} - E^{V} \|_{F}^{2} + \tau \| \alpha \|_{1}$$
(10)

$$\min_{\alpha} \| E^{U}C - C \sum_{i} \alpha_{i} E^{V_{i}} \|^{2} + \tau \| \alpha \|_{1}$$
(14)

$$E(\widetilde{\mathsf{L}}) = \|\mathsf{P}\mathsf{M}^{-1}\mathsf{L}\mathsf{I} - \widetilde{\mathsf{M}}^{-1}\widetilde{\mathsf{L}}\mathsf{P}\mathsf{I}\|_{\widetilde{\mathsf{M}}}^{2},$$
(3)

$$\min_{\{\mathbf{a}_l,\mathbf{b}_l,\mathbf{c}_l\}} \|R - \sum_{l=1}^m \mathbf{a}_l \otimes \mathbf{b}_l \otimes \mathbf{c}_l\|_{\mathcal{F}}^2$$

$$\gamma \sum_{i=1}^{r} A_1(p_i) \|X_{12}(p_i,:) - X_2(q_i)\|_{M_2}^2 + A_2(q_i) \|X_{21}(q_i,:) - X_1(p_i)\|_{M_2}^2$$

$$\mathbf{w}_{j} = \left(\mathbf{R}_{j}^{W}\right)^{\top} \begin{pmatrix} \mathbf{P} - \mathbf{t}_{j}^{W} \\ \|\mathbf{P} - \mathbf{t}_{j}^{W}\|_{2} \end{pmatrix}.$$
 (13)





Norms in I LA

 $a = ||T||_1 + ||T||$ $b = ||T||_{\infty} + ||T||_P$ $c = ||P||_* + ||P||_F$ where $T: \mathbb{R}^2$: a vector $P: \mathbb{R}^{(2\times 2)}$: a matrix



















myExpressionResultType myExpression(const Eigen::Matrix<double, 2, 1> & x, const std::function<double(Eigen::Matrix<double, 2, 1>)> & f)

double y = pow(f(x), 2);
return myExpressionResultType(y);

C++







myExpressionResultType myExpression(
 const Eigen::Matrix<double, 2, 1> & x,
 const std::function<double(Eigen::Matrix<double, 2, 1>)> & f)

double y = pow(f(x), 2);
return myExpressionResultType(y);

def myExpression(x, f): x = np.asarray(x, dtype=np.float64) assert x.shape == (2,) y = np.power(f(x), 2) return myExpressionResultType(y)







myExpressionResultType myExpression(
 const Eigen::Matrix<double, 2, 1> & x,
 const std::function<double(Eigen::Matrix<double, 2, 1>)> & f)

double y = pow(f(x), 2);
return myExpressionResultType(y);

def myExpression(x, f): x = np.asarray(x, dtype=np.float64) assert x.shape == (2,) y = np.power(f(x), 2) return myExpressionResultType(y)

```
function output = myExpression(x, f)
    x = reshape(x,[],1);
    assert( numel(x) == 2 );
    y = f(x).^2;
    output.y = y;
end
```









Generation

end{align*}













\end{align*}

. . .











(z) Geometry Processing Course: Registration [Jacobson 2020]

(w) Plenoptic Modeling: An Image-Based Rendering System [McMillan and Bishop 1995] Eq. 22

$$L_{ij} = \begin{cases} w_{ij} \\ -\sum_{\ell \neq i} L_{i\ell} \\ 0 \end{cases}$$

- $L_i, j = \{ w_i, j if (i, j) \in E \}$ **0** otherwise
- $L_i, i = -\sum_{\ell \in I} (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ w E R^(n×n): edge weight matrix $E \in \{\mathbb{Z}^2\}$ index: edges

(a) Geometry Processing Course: Parameterization [Jacobson 2020]

Examples from the Wild









$$L_{ij} = \begin{cases} w_{ij} \\ -\sum_{\ell \neq i} L_{i\ell} \\ 0 \end{cases}$$

 $L_i,j = \{ w_i,j if (i,j) \in E \}$ **0** otherwise $L_i, i = -\sum_{\ell \in I} (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ $w \in \mathbb{R}^{(n \times n)}$: edge weight matrix $E \in \{\mathbb{Z}^2\}$ index: edges

Examples from the Wild

if $i \neq j$ and $\exists \{ij\} \in \mathbf{E}$ if i = j, or

otherwise



L_i,j = { w_i,j if (i,j)
$$\in E$$

0 otherwise
L_i,i = $-\sum_{\ell} (\ell \text{ for } \ell \neq i) L_i$,
where
L $\in \mathbb{R}^{n\times n}$
w $\in \mathbb{R}^{n\times n}$: edge weight matr
E $\in \{\mathbb{Z}^2\}$ index: edges



matrix

An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\}\in \mathbf{E}} w_{ij} \|\mathbf{u}_i - \mathbf{u}_j\|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame *length distortion*. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage *area distortion* and *angle distortion*.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n imes n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E}, \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}$$

What's up with the $\operatorname{tr}()$ in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the trace of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to \mathbf{U} . In effect, this is really applying \mathbf{Q} to each column of \mathbf{U} independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

I\U0076LA implementation:

L_i,j = { w_i,j if (i,j) ∈ E Ø otherwise L_i,i = -∑_(ℓ for ℓ ≠ i) L_i,ℓ

where

L ∈ ℝ^(n×n)

- w $\in \mathbb{R}^{(n \times n)}$: edge weight matrix
- $E \in \{\mathbb{Z}^2\}$ index: edges



An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\} \in \mathbf{E}} w_{ij} \|\mathbf{u}_i - \mathbf{u}_j\|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame *length distortion*. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage *area distortion* and *angle distortion*.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n imes n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E}, \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}$$

What's up with the ${\rm tr}\left(\right)$ in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the trace of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to \mathbf{U} . In effect, this is really applying \mathbf{Q} to each column of \mathbf{U} independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

I\U0076LA compiled to C++/Eigen:

```
*/
course_parameterizationResultType course_parameterization(
   const Eigen::MatrixXd & w,
   const std::set<std::tuple< int, int > > & E)
{
   const long n = w.cols();
   assert( w.rows() == n );
   Eigen::SparseMatrix<double> L(n, n);
   std::vector<Eigen::Triplet<double> > tripletList_L;
   for( int i=1; i<=n; i++){</pre>
        for( int j=1; j<=n; j++){</pre>
            if(E.find(std::tuple< int, int >(i-1, j-1)) != E.end()){
                tripletList_L.push_back(Eigen::Triplet<double>(i-1, j-1)
   L.setFromTriplets(tripletList_L.begin(), tripletList_L.end());
   for( int i=1; i<=n; i++){</pre>
        double sum_0 = 0;
```

I\U0076LA implementation:

L_i,j = { w_i,j if (i,j) ∈ E 0 otherwise L_i,i = -∑_(ℓ for ℓ ≠ i) L_i,ℓ

where $L \in \mathbb{R}^{n\times n}$

- w $\in \mathbb{R}^{(n \times n)}$: edge weight matrix
- $E \in \{\mathbb{Z}^2\}$ index: edges



An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\} \in \mathbf{E}} w_{ij} \| \mathbf{u}_i - \mathbf{u}_j \|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame length distortion. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage area distortion and angle distortion.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n \times n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E}, \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}.$$

What's up with the tr() in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the trace of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to U. In effect, this is really applying \mathbf{Q} to each column of \mathbf{U} independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

I VLA compiled to C++/Eigen:

```
*/
course_parameterizationResultType course_parameterization(
    const Eigen::MatrixXd & w,
    const std::set<std::tuple< int, int > > & E)
    const long n = w.cols();
    assert( w.rows() == n );
    Eigen::SparseMatrix<double> L(n, n);
    std::vector<Eigen::Triplet<double> > tripletList_L;
    for( int i=1; i<=n; i++){</pre>
        for( int j=1; j<=n; j++){</pre>
            if(E.find(std::tuple< int, int >(i-1, j-1)) != E.end()){
                tripletList_L.push_back(Eigen::Triplet<double>(i-1, j-1)
    L.setFromTriplets(tripletList_L.begin(), tripletList_L.end());
    for( int i=1; i<=n; i++){</pre>
        double sum_0 = 0;
```

I*LA implementation:

 $L_i, j = \{ w_i, j \text{ if } (i, j) \in E \}$ 0 otherwise $L_i, i = -\sum_{\ell} (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ w $\in \mathbb{R}^{(n \times n)}$: edge weight matrix

 $E \in \{\mathbb{Z}^2\}$ index: edges

```
LA compiled to Python/NumPy/SciPy
def course_parameterization(w, E):
    .....
    :param :w : edge weight matrix
    :param :E : edges
    ......
    w = np.asarray(w, dtype=np.float64)
    E = frozenset(E)
    n = w.shape[1]
    assert w.shape == (n, n)
    Lij_0 = []
    Lvals_0 = 🗌
    for i in range(1, n+1):
        for j in range(1, n+1):
            if (i-1, j-1) in E:
    L = sparse_0
    for i in range(1, n+1):
```



:		

 $Lij_0.append((i-1, j-1))$ Lvals_0.append(w[i-1, j-1]) sparse_0 = scipy.sparse.coo_matrix((Lvals_0, np.asarray(Lij_0).T),

An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\} \in \mathbf{E}} w_{ij} \| \mathbf{u}_i - \mathbf{u}_j \|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame *length distortion*. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage *area distortion* and *angle distortion*.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n imes n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E}, \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}.$$

What's up with the ${\rm tr}\left(\right)$ in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the **trace** of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to \mathbf{U} . In effect, this is really applying \mathbf{Q} to each column of \mathbf{U} independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

WLA compiled to C++/Eigen:

```
*/
course_parameterizationResultType course_parameterization(
    const Eigen::MatrixXd & w,
    const std::set<std::tuple< int, int >> & E)
{
    const long n = w.cols();
    assert( w.rows() == n );
    Eigen::SparseMatrix<double> L(n, n);
    std::vector<Eigen::Triplet<double> > tripletList_L;
    for( int i=1; i<=n; i++){
        for( int j=1; j<=n; j++){
            if(E.find(std::tuple< int, int >(i-1, j-1)) != E.end()){
            tripletList_L.push_back(Eigen::Triplet<double>(i-1, j=1)
            }
        }
        L.setFromTriplets(tripletList_L.begin(), tripletList_L.end());
        for( int i=1; i<=n; i++){
            double sum_0 = 0;
        }
    }
}
</pre>
```

I*LA implementation:

 $L_i, j = \{ w_i, j \text{ if } (i, j) \in E \\ 0 \text{ otherwise} \\ L_i, i = -\sum_{\ell} (\ell \text{ for } \ell \neq i) L_i, \ell \}$

where
L ∈ ℝ^(n×n)
w ∈ ℝ^(n×n): edge weight matrix

 $E \in \{\mathbb{Z}^2\}$ index: edges

```
LA compiled to Python/NumPy/SciPy:
def course_parameterization
    :param :w : edge weight
    :param :E : edges
    .....
    w = np.asarray(w, dtype)
    E = frozenset(E)
    n = w.shape[1]
    assert w.shape == (n, n)
    Lij_0 = []
    Lvals_0 = 📋
    for i in range(1, n+1):
        for j in range(1, n
            if (i-1, j-1) i
                Lij_0.appen
                Lvals_0.app
    sparse_0 = scipy.sparse
   L = sparse_0
    for i in range(1, n+1):
```



(w, E):
matrix
=np.float64)
)
+1): n E: d((i-1, j-1)) end(w[i-1, j-1]) .coo_matrix((Lvals_0, np.asarray(Lij_0).T),

LA compiled to MATLAB:

```
Lvals_0 = zeros(1,0);
for i = 1:n
    for j = 1:n
       if ismember([i, j],E,'rows')
            Lij_0(1:2, end+1) = [i;j];
            Lvals_0(end+1) = w(i, j);
        end
    end
end
sparse_0 = sparse(Lij_0(1,:),Lij_0(2,:),Lvals_0,n,n);
L = sparse_0;
for i = 1:n
    sum_0 = 0;
    for ell = 1:size(L,2)
        if ell ~= i
          sum_0 = sum_0 + L(i, ell);
        end
    end
    Lij_0(1:2, end+1) = [i;i];
```



An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\} \in \mathbf{E}} w_{ij} \| \mathbf{u}_i - \mathbf{u}_j \|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame length distortion. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage area distortion and angle distortion.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n \times n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E}, \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}.$$

What's up with the tr() in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the trace of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to U. In effect, this is really applying Q to each column of U independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

LA compiled to C++/Eigen:

```
*/
course_parameterizationResultType course_parameterization(
    const Eigen::MatrixXd & w,
    const std::set<std::tuple< int, int > > & E)
    const long n = w.cols();
    assert( w.rows() == n );
    Eigen::SparseMatrix<double> L(n, n);
    std::vector<Eigen::Triplet<double> > tripletList_L;
    for( int i=1; i<=n; i++){</pre>
        for( int j=1; j<=n; j++){</pre>
            if(E.find(std::tuple< int, int >(i-1, j-1)) != E.end()){
                tripletList_L.push_back(Eigen::Triplet<double>(i-1, j-1)
    L.setFromTriplets(tripletList_L.begin(), tripletList_L.end());
    for( int i=1; i<=n; i++){</pre>
        double sum_0 = 0;
```

I*LA implementation:

 $L_{i,j} = \{ w_{i,j} \text{ if } (i,j) \in E \}$ 0 otherwise $L_i, i = -\sum (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ w $\in \mathbb{R}^{(n \times n)}$: edge weight matrix

 $E \in \{\mathbb{Z}^2\}$ index: edges

```
LA compiled to Python/NumPy/SciPy:
def course_parameterization(w, E):
    :param :w : edge weight matrix
    :param :E : edges
    w = np.asarray(w, dtype=np.float64)
    E = frozenset(E)
    n = w.shape[1]
    assert w.shape == (n, n)
    Lij_0 = []
    Lvals_0 = 🗌
    for i in range(1, n+1):
        for j in range(1, n+1):
            if (i-1, j-1) in E:
   L = sparse_0
    for i in range(1, n+1):
```



 $Lij_0.append((i-1, j-1))$ Lvals_0.append(w[i-1, j-1]) sparse_0 = scipy.sparse.coo_matrix((Lvals_0, np.asarray(Lij_0).T),

LA compiled to MATLAB:

```
Lvals_0 = zeros(1,0);
for i = 1:n
    for j = 1:n
       if ismember([i, j],E,'rows')
            Lij_0(1:2, end+1) = [i;j];
            Lvals_0(end+1) = w(i, j);
        end
    end
end
sparse_0 = sparse(Lij_0(1,:),Lij_0(2,:),Lvals_0,n,n);
L = sparse_0;
for i = 1:n
    sum_0 = 0;
    for ell = 1:size(L,2)
        if ell ~= i
          sum_0 = sum_0 + L(i, ell);
        end
    end
    Lij_0(1:2, end+1) = [i;i];
```



An example from Geometry Processing: Parameterization

The original equation:

We can try to remedy this by introducing a non-uniform weight or spring stiffness w_{ij} for each edge $\{i, j\}$:

$$\min_{\mathbf{U}} \sum_{\{i,j\} \in \mathbf{E}} w_{ij} \| \mathbf{u}_i - \mathbf{u}_j \|^2.$$

For example, we could weigh the distortion of shorter edges (on the 3D mesh) more than longer ones: $w_{ij} = 1/||\mathbf{v}_i - \mathbf{v}_j||$. See "Parametrization and smooth approximation of surface triangulations" [Floater 1996]. This will at best help tame length distortion. The "shapes" (i.e., aspect ratios) of triangles will only be indirectly preserved. We need a way to discourage area distortion and angle distortion.

To do this, let's write the energy minimization problem above in matrix form:

$$\min_{\mathbf{U}} \frac{1}{2} \operatorname{tr} \big(\mathbf{U}^{\mathsf{T}} \mathbf{L} \mathbf{U} \big),$$

where $\mathbf{L} \in \mathbb{R}^{n \times n}$ is a sparse matrix with:

$$L_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \text{ and } \exists \{ij\} \in \mathbf{E} \\ -\sum_{\ell \neq i} L_{i\ell} & \text{if } i = j, \text{ or} \\ 0 & \text{otherwise} \end{cases}$$

What's up with the tr() in the energy?

The degrees of freedom in our optimization are a collected in the *matrix* $\mathbf{U} \in \mathbb{R}^{n \times 2}$ with two columns. The energy is written as the trace of the quadratic form (a.k.a. matrix) $\mathbf{Q} \in \mathbb{R}^{n \times n}$ applied to U. In effect, this is really applying \mathbf{Q} to each column of \mathbf{U} independently and summing the result:

$$\begin{aligned} \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) &= \\ &= \operatorname{tr} \left(\mathbf{U}^{\mathsf{T}} \mathbf{Q} \mathbf{U} \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \\ \mathbf{U}_{2}^{\mathsf{T}} \end{bmatrix} \mathbf{Q} [\mathbf{U}_{1} \mathbf{U}_{2}] \right) \\ &= \operatorname{tr} \left(\begin{bmatrix} \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \\ \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} & \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2} \end{bmatrix} \\ &= \mathbf{U}_{1}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{1} + \mathbf{U}_{2}^{\mathsf{T}} \mathbf{Q} \mathbf{U}_{2}. \end{aligned}$$

WLA compiled to C++/Eigen:

```
*/
course_parameterizationResultType course_parameterization(
    const Eigen::MatrixXd & w,
    const std::set<std::tuple< int, int > > & E)
    const long n = w.cols();
    assert( w.rows() == n );
    Eigen::SparseMatrix<double> L(n, n);
    std::vector<Eigen::Triplet<double> > tripletList_L;
    for( int i=1; i<=n; i++){</pre>
        for( int j=1; j<=n; j++){</pre>
            if(E.find(std::tuple< int, int >(i-1, j-1)) != E.end()){
                tripletList_L.push_back(Eigen::Triplet<double>(i-1, j-1)
    L.setFromTriplets(tripletList_L.begin(), tripletList_L.end());
    for( int i=1; i<=n; i++){</pre>
        double sum_0 = 0;
```

I*LA implementation:

$$L_i, j = \{ w_i, j \text{ if } (i, j) \\ 0 \text{ otherwise} \\ L_i, i = -\sum_{\ell \in I} (\ell \text{ for } \ell \neq i) \}$$

where

- $L \in \mathbb{R}^{(n \times n)}$
- w E R^(n×n): edge weight matrix
- **E** $\{\mathbb{Z}^2\}$ index: edges

def course_parameterization(w, E): :param :w : edge weight matrix :param :E : edges w = np.asarray(w, dtype=np.float64)E = frozenset(E)n = w.shape[1]assert w.shape == (n, n) Lij_0 = [] Lvals_0 = 🗌 for i in range(1, n+1): for j in range(1, n+1):

if (i-1, j-1) in E: $L = sparse_0$ for i in range(1, n+1):



LA compiled to Python/NumPy/SciPy:

 $Lij_0.append((i-1, j-1))$ Lvals_0.append(w[i-1, j-1]) sparse_0 = scipy.sparse.coo_matrix((Lvals_0, np.asarray(Lij_0).T),

LA compiled to MATLAB:

```
Lvals_0 = zeros(1,0);
for i = 1:n
    for j = 1:n
       if ismember([i, j],E,'rows')
            Lij_0(1:2, end+1) = [i;j];
           Lvals_0(end+1) = w(i, j);
        end
    end
end
sparse_0 = sparse(Lij_0(1,:),Lij_0(2,:),Lvals_0,n,n);
L = sparse_0;
for i = 1:n
    sum_0 = 0;
    for ell = 1:size(L,2)
        if ell ~= i
          sum_0 = sum_0 + L(i, ell);
        end
    end
    Lij_0(1:2, end+1) = [i;i];
```



$$L_{ij} = \begin{cases} w_{ij} \\ -\sum_{\ell \neq i} L_{i\ell} \\ 0 \end{cases}$$

 $L_i,j = \{ w_i,j if (i,j) \in E \}$ **0** otherwise $L_i, i = -\sum_{\ell \in I} (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ $w \in \mathbb{R}^{(n \times n)}$: edge weight matrix $E \in \{\mathbb{Z}^2\}$ index: edges

Examples from the Wild

if $i \neq j$ and $\exists \{ij\} \in \mathbf{E}$ if i = j, or

otherwise

$$L_{ij} = \begin{cases} w_{ij} \\ -\sum_{\ell \neq i} L_{i\ell} \\ 0 \end{cases}$$

- $L_i, j = \{ w_i, j if (i, j) \in E \}$ **0** otherwise
- $L_i, i = -\sum_{\ell \in I} (\ell \text{ for } \ell \neq i) L_i, \ell$

where $L \in \mathbb{R}^{(n \times n)}$ w E R^(n×n): edge weight matrix $E \in \{\mathbb{Z}^2\}$ index: edges

(a) Geometry Processing Course: Parameterization [Jacobson 2020]

Examples from the Wild











(z) Geometry Processing Course: Registration [Jacobson 2020]

(w) Plenoptic Modeling: An Image-Based Rendering System [McMillan and Bishop 1995] Eq. 22
- **Polygon Mesh Processing Library**
- **Conforming Weighted Delaunay Triangulations**
- Properties of Laplace Operators for Tetrahedral Meshes_volume
- **Properties of Laplace Operators for Tetrahedral Meshes_circumcenter**
- **Instant Field-Aligned Meshes**
- Collision-Aware and Online Compression of Rigid Body Simulations via Integrated Error Minimization
- Frame Fields: Anisotropic and Non-Orthogonal Cross Fields
- **Regularized Kelvinlets: Sculpting Brushes based on Fundamental Solutions of Elasticity**
- Robust Inside-Outside Segmentation using Generalized Winding Numbers
- **Gaussian-Product Subdivision Surfaces**

- **Polygon Mesh Processing Library**
- **Conforming Weighted Delaunay Triangulations**
- Properties of Laplace Operators for Tetrahedral Meshes_volume
- **Properties of Laplace Operators for Tetrahedral Meshes_circumcenter**
- **Instant Field-Aligned Meshes**
- Collision-Aware and Online Compression of Rigid Body Simulations via Integrated Error Minimization
- Frame Fields: Anisotropic and Non-Orthogonal Cross Fields
- **Regularized Kelvinlets: Sculpting Brushes based on Fundamental Solutions of Elasticity**
- Robust Inside-Outside Segmentation using Generalized Winding Numbers
- **Gaussian-Product Subdivision Surfaces**











- **Polygon Mesh Processing Library**
- **Conforming Weighted Delaunay Triangulations**
- Properties of Laplace Operators for Tetrahedral Meshes_volume
- **Properties of Laplace Operators for Tetrahedral Meshes_circumcenter**
- **Instant Field-Aligned Meshes**
- Collision-Aware and Online Compression of Rigid Body Simulations via Integrated Error Minimization
- Frame Fields: Anisotropic and Non-Orthogonal Cross Fields
- Regularized Kelvinlets: Sculpting Brushes based on Fundamental Solutions of Elasticity
- Robust Inside-Outside Segmentation using Generalized Winding Numbers
- **Gaussian-Product Subdivision Surfaces**











Integrating with existing code from Instant Field-Aligned Meshes.

Project URL: Instant Field-Aligned Meshes

The original formula:

Intermediate position. We define a position q_{ij} that minimizes the distance to vertices v_i and v_j while being located in their respective tangent planes, i.e.:

minimize
$$\|\mathbf{v}_i - \mathbf{q}_{ij}\|_2^2 + \|\mathbf{v}_j - \mathbf{q}_{ij}\|_2^2$$

subject to $\langle \mathbf{n}_i, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_i, \mathbf{v}_i \rangle$ and $\langle \mathbf{n}_j, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_j, \mathbf{v}_j \rangle$.

This constrained least-squares problem has a simple solution:

$$\mathbf{q}_{ij} = \frac{1}{2} (\mathbf{v}_i + \mathbf{v}_j) - \frac{1}{4} (\lambda_i \mathbf{n}_i + \lambda_j \mathbf{n}_j),$$

where the Lagrange multiplier λ_i is

$$\lambda_i = \frac{2\left\langle (\mathbf{n}_i + \langle \mathbf{n}_i, \mathbf{n}_j \rangle \mathbf{n}_j, \, \mathbf{v}_j - \mathbf{v}_i \right\rangle}{1 - \langle \mathbf{n}_i, \mathbf{n}_j \rangle^2 + \varepsilon},$$

and λ_j is defined analogously with *i* and *j* swapped. The parameter ε (set to 10^{-4} in our implementation) ensures that \mathbf{q}_{ij} approximates the arithmetic mean of \mathbf{v}_i and \mathbf{v}_j when $\mathbf{n}_i \approx \mathbf{n}_j$.

Integrating with existing code from Instant Field-Aligned Meshes.

Project URL: Instant Field-Aligned Meshes

The original formula:

Intermediate position. We define a position q_{ij} that minimizes the distance to vertices v_i and v_j while being located in their respective tangent planes, i.e.:

minimize
$$\|\mathbf{v}_i - \mathbf{q}_{ij}\|_2^2 + \|\mathbf{v}_j - \mathbf{q}_{ij}\|_2^2$$

subject to $\langle \mathbf{n}_i, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_i, \mathbf{v}_i \rangle$ and $\langle \mathbf{n}_j, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_j, \mathbf{v}_j \rangle$.

This constrained least-squares problem has a simple solution:

$$\mathbf{q}_{ij} = \frac{1}{2} (\mathbf{v}_i + \mathbf{v}_j) - \frac{1}{4} (\lambda_i \mathbf{n}_i + \lambda_j \mathbf{n}_j),$$

where the Lagrange multiplier λ_i is

$$\lambda_i = \frac{2\left\langle (\mathbf{n}_i + \langle \mathbf{n}_i, \mathbf{n}_j \rangle \mathbf{n}_j, \, \mathbf{v}_j - \mathbf{v}_i \right\rangle}{1 - \langle \mathbf{n}_i, \mathbf{n}_j \rangle^2 + \varepsilon},$$

The <u>original source code</u>:

		o- field.cpp — instant-meshes
FOLDERS	4 •	field.cpp ×
🔻 📄 instant-meshes	0 8	return q;
▶ 🖿 build	8.	
▶ ■ ext	84	inline Vector3f middle_point(const Vector3f &p0, const Vector3f &n0, const Vector3f &p1, const Vector3f &n1) {
	8	/* How was this derived?
	8	*
src src	8	$\frac{1}{3} + \frac{1}{10} +$
/* aabb.h	8	* dot(n1, x) == dot(n1, p1)
/* adjacency.cpp	90) *
/* adjacency.h	91	* -> Lagrange multipliers, set derivative = 0
/* batch cpp	9.	* Use first 3 equalifies to write X in terms of * lambda 1 and lambda 2. Substitute that into the last
	94	* two equations and solve for the lambdas. Finally.
/* batch.h	9!	* add a small epsilon term to avoid issues when $n1=n2$.
/* bvh.cpp	91	i */
∕∗ bvh.h	9	Float $n0p0 = n0.dot(p0), n0p1 = n0.dot(p1),$
/* cleanup.cpp	91	n1p0 = n1.dot(p0), n1p1 = n1.dot(p1),
	10	denom = $1.0f / (1.0f - n0n1*n0n1 + 1e-4f)$.
/* cleanup.n	10	lambda_0 = 2.0f*(n0p1 - n0p0 - n0n1*(n1p0 - n1p1))*denom,
/* common.h	10	<pre>lambda_1 = 2.0f*(n1p0 - n1p1 - n0n1*(n0p1 - n0p0))*denom;</pre>
∕∗ dedge.cpp	10	
∕∗ dedge.h	104 105	<pre>return 0.5f * (p0 + p1) - 0.25f * (n0 * lambda_0 + n1 * lambda_1); }</pre>
∕∗ diff.cpp	10	
/* extract.cpp	10	<pre>std::pair<vector3f, vector3f=""> compat_orientation_intrinsic_2(const Vector3f & a0, const Vector3f & a1, const Vector3f & a1) {</vector3f,></pre>
/* extract.h	10	const Vector3f q1 = rotate_vector_into_plane(_q1, n1, n0);
/* field.cpp	110 111	<pre>preturn std::make_pair(q0, q1 * signum(q1.dot(q0))); }</pre>



Integrating with existing code from Instant Field-Aligned Meshes.

Project URL: Instant Field-Aligned Meshes

The original formula:

Intermediate position. We define a position q_{ij} that minimizes the distance to vertices v_i and v_j while being located in their respective tangent planes, i.e.:

minimize
$$\|\mathbf{v}_i - \mathbf{q}_{ij}\|_2^2 + \|\mathbf{v}_j - \mathbf{q}_{ij}\|_2^2$$

subject to $\langle \mathbf{n}_i, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_i, \mathbf{v}_i \rangle$ and $\langle \mathbf{n}_j, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_j, \mathbf{v}_j \rangle$.

This constrained least-squares problem has a simple solution:

$$\mathbf{q}_{ij} = \frac{1}{2} (\mathbf{v}_i + \mathbf{v}_j) - \frac{1}{4} (\lambda_i \mathbf{n}_i + \lambda_j \mathbf{n}_j),$$

where the Lagrange multiplier λ_i is

$$\lambda_i = \frac{2\left\langle (\mathbf{n}_i + \langle \mathbf{n}_i, \mathbf{n}_j \rangle \mathbf{n}_j, \ \mathbf{v}_j - \mathbf{v}_i \right\rangle}{1 - \langle \mathbf{n}_i, \mathbf{n}_j \rangle^2 + \varepsilon},$$

The original source code:

		o field.cpp — instant-meshes
FOLDERS	•	field.cpp ×
🔻 📄 instant-meshes	0	81 return q;
▶ 🛄 build		82 }
▶ 📰 ext		84 inline Vector3f middle_point(const Vector3f &p0, const Vector3f &n0, const Vector3f &p1, const Vector3f &n1) {
resources		85 /* How was this derived?
	0	$\frac{86}{97} \times \frac{1}{2} \times $
src	0	88 + dot(n0, x) == dot(n0, n0)
∕∗ aabb.h		89 + dot(n1, x) == dot(n1, p1)
/* adjacency.cpp		90 *
/* adjacency h		91 * -> Lagrange multipliers, set derivative = 0
		92 * Use first 3 equalities to write x in terms of
/* batch.cpp		93 * lambda_1 and lambda_2. Substitute that into the last
/* batch.h		94 * two equations and solve for the lambdas. Finally,
/* byh.cpp		95 * add a small epsilon term to avoid issues when ni=n2.
		97 Float $n0n0 = n0.dot(n0), n0n1 = n0.dot(n1),$
/* bvn.n		98 $n1p0 = n1.dot(p0), n1p1 = n1.dot(p1),$
/* cleanup.cpp		99 $n0n1 = n0.dot(n1),$
/* cleanup.h		denom = 1.0f / (1.0f - n0n1*n0n1 + 1e-4f),
		01 lambda_0 = 2.0f*(n0p1 - n0p0 - n0n1*(n1p0 - n1p1))*denom,
/* common.n		02 lambda_1 = 2.0f*(n1p0 - n1p1 - n0n1*(n0p1 - n0p0))*denom;
/* dedge.cpp		
∕∗ dedge.h		04
∕∗ diff.cpp		06
/* extract.cpp		07 std::pair <vector3f, vector3f=""> compat_orientation_intrinsic_2(</vector3f,>
/* extract.h		09 const Vector3f q1 = rotate_vector_into_plane(_q1, n1, n0);
/* field.cpp		<pre>10 return std::make_pair(q0, q1 * signum(q1.dot(q0))); 11 }</pre>

$$\begin{aligned} & \epsilon = 10^{(-4)} \\ \hat{\lambda}_{i} &= (2(\hat{n}_{i} + (\hat{n}_{i}, \hat{n}_{j})) \hat{n}_{j}, \hat{v}_{j} - \hat{v}_{i}))/(1 - (\hat{n}_{i}, \hat{n}_{j})^{2} + \epsilon) \\ \hat{\lambda}_{j} &= (2(\hat{n}_{j} + (\hat{n}_{j}), \hat{n}_{i})) \hat{n}_{i}, \hat{v}_{i} - \hat{v}_{j})/(1 - (\hat{n}_{j}), \hat{n}_{i})^{2} + \epsilon) \\ \hat{q}_{ij} &= 1/2(\hat{v}_{i} + \hat{v}_{j}) - 1/4(\hat{\lambda}_{i}) \hat{n}_{i} + \hat{\lambda}_{j}) \hat{n}_{j}) \end{aligned}$$

where

`

`vi` ∈ ℝ^3 `ni` ∈ ℝ^3 `v;` ∈ ℝ^3 `n;` ∈ ℝ^3





Integrating with existing code from Instant Field-Aligned Meshes.

Project URL: Instant Field-Aligned Meshes

The original formula:

Intermediate position. We define a position q_{ij} that minimizes the distance to vertices \mathbf{v}_i and \mathbf{v}_j while being located in their respective tangent planes, i.e.:

minimize
$$\|\mathbf{v}_i - \mathbf{q}_{ij}\|_2^2 + \|\mathbf{v}_j - \mathbf{q}_{ij}\|_2^2$$

subject to $\langle \mathbf{n}_i, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_i, \mathbf{v}_i \rangle$ and $\langle \mathbf{n}_j, \mathbf{q}_{ij} \rangle = \langle \mathbf{n}_j, \mathbf{v}_j \rangle$.

This constrained least-squares problem has a simple solution:

$$\mathbf{q}_{ij} = \frac{1}{2} (\mathbf{v}_i + \mathbf{v}_j) - \frac{1}{4} (\lambda_i \mathbf{n}_i + \lambda_j \mathbf{n}_j),$$

where the Lagrange multiplier λ_i is

$$\lambda_i = \frac{2\left\langle (\mathbf{n}_i + \langle \mathbf{n}_i, \mathbf{n}_j \rangle \mathbf{n}_j, \ \mathbf{v}_j - \mathbf{v}_i \right\rangle}{1 - \langle \mathbf{n}_i, \mathbf{n}_j \rangle^2 + \varepsilon},$$

The <u>original source code</u>:

		o field.cpp — instant-meshes
FOLDERS	•	field.cpp ×
🔻 📄 instant-meshes	0	81 return q;
▶ 🛄 build		82 }
▶ 📰 ext		84 inline Vector3f middle_point(const Vector3f &p0, const Vector3f &n0, const Vector3f &p1, const Vector3f &n1) {
resources		85 /* How was this derived?
	0	$\frac{86}{97} \times \frac{1}{2} \times $
src	0	88 + dot(n0, x) == dot(n0, n0)
∕∗ aabb.h		89 + dot(n1, x) == dot(n1, p1)
/* adjacency.cpp		90 *
/* adjacency h		91 * -> Lagrange multipliers, set derivative = 0
		92 * Use first 3 equalities to write x in terms of
/* batch.cpp		93 * lambda_1 and lambda_2. Substitute that into the last
/* batch.h		94 * two equations and solve for the lambdas. Finally,
/* byh.cpp		95 * add a small epsilon term to avoid issues when ni=n2.
		97 Float $n0n0 = n0.dot(n0), n0n1 = n0.dot(n1),$
/* bvn.n		98 $n1p0 = n1.dot(p0), n1p1 = n1.dot(p1),$
/* cleanup.cpp		99 $n0n1 = n0.dot(n1),$
/* cleanup.h		denom = 1.0f / (1.0f - n0n1*n0n1 + 1e-4f),
		01 lambda_0 = 2.0f*(n0p1 - n0p0 - n0n1*(n1p0 - n1p1))*denom,
/* common.n		02 lambda_1 = 2.0f*(n1p0 - n1p1 - n0n1*(n0p1 - n0p0))*denom;
/* dedge.cpp		
∕∗ dedge.h		04
∕∗ diff.cpp		06
/* extract.cpp		07 std::pair <vector3f, vector3f=""> compat_orientation_intrinsic_2(</vector3f,>
/* extract.h		09 const Vector3f q1 = rotate_vector_into_plane(_q1, n1, n0);
/* field.cpp		<pre>10 return std::make_pair(q0, q1 * signum(q1.dot(q0))); 11 }</pre>

$$= 10^{(-4)}$$
 $\lambda_i = (2(n_i + (n_i), n_j))n_j, v_j - v_i))/(1 - (n_i), n_j)^2 + \epsilon)$
 $\lambda_j = (2(n_j + (n_j), n_i))n_i, v_i - v_j))/(1 - (n_j), n_i)^2 + \epsilon)$
 $q_{ij} = 1/2(v_i + v_j) - 1/4(\lambda_i)n_i + \lambda_j)n_j)$

E



instant-meshes-SA-2015-jakob-et-al E ~ Page 1 of 15

Q Ð

Instant_iField-Aligned Meshes

Wenzel Jakob¹

Marco Tarini^{2,3}

¹ETH Zurich ²CNR-ISTI



Figure 1: Remeshing a scanned dragon with 13 million vertices into feature-aligned isotropic triangle and quad meshes with \sim 80k vertices. From left to right, for both cases: visualizations of the orientation field, position field, and the output mesh (computed in 71.1 and 67.2 seconds, respectively). For the quad case, we optimize for a quad-dominant mesh at quarter resolution and subdivide once to obtain a pure quad mesh.

Abstract

We present a novel approach to remesh a surface into an isotropic triangular or quad-dominant mesh using a unified local smoothing operator that optimizes both the edge orientations and vertex positions in the output mesh. Our algorithm produces meshes with high isotropy while naturally aligning and snapping edges to sharp features. The method is simple to implement and parallelize, and it can process a variety of input surface representations, such as point clouds, range scans and triangle meshes. Our full pipeline executes instantly (less than a second) on meshes with hundreds of thousands of faces, enabling new types of interactive workflows. Since our algorithm avoids any global optimization, and its key steps scale linearly with input size, we are able to process extremely large meshes and point clouds, with sizes exceeding several hundred INSERT MODE, 12 characters selected



INSERT MODE, Line 9, Column 1



Daniele Panozzo¹

Olga Sorkine-Hornung¹

³Università dell'Insubria

Introduction

Triangle and quad-dominant meshes are ubiquitously used in computer graphics and CAD applications to represent surfaces, either directly, or as the control grid for higher-order parametric or subdivision surfaces. With the introduction of T-splines [Sederberg et al. 2003] and Dyadic T-Mesh Subdivision [Kovacs et al. 2015], quad-dominant meshes with T-joints (T-meshes) now have similar properties and applications of pure quadrilateral meshes, while being more flexible and naturally supporting the flexible local refinement that is often desired in CAD applications. Meshing surfaces is a challenging problem, and a plethora of methods have been proposed in the past three decades to cope with the increasing quality and scalability requirements of modern applications [Owen 1998; Bommes et al. 2013a]. Semi-regular meshes, which have uniform element



instant-meshes-SA-2015-jakob-et-al E ~ Page 1 of 15

Q Ð

Instant_iField-Aligned Meshes

Wenzel Jakob¹

Marco Tarini^{2,3}

¹ETH Zurich ²CNR-ISTI



Figure 1: Remeshing a scanned dragon with 13 million vertices into feature-aligned isotropic triangle and quad meshes with \sim 80k vertices. From left to right, for both cases: visualizations of the orientation field, position field, and the output mesh (computed in 71.1 and 67.2 seconds, respectively). For the quad case, we optimize for a quad-dominant mesh at quarter resolution and subdivide once to obtain a pure quad mesh.

Abstract

We present a novel approach to remesh a surface into an isotropic triangular or quad-dominant mesh using a unified local smoothing operator that optimizes both the edge orientations and vertex positions in the output mesh. Our algorithm produces meshes with high isotropy while naturally aligning and snapping edges to sharp features. The method is simple to implement and parallelize, and it can process a variety of input surface representations, such as point clouds, range scans and triangle meshes. Our full pipeline executes instantly (less than a second) on meshes with hundreds of thousands of faces, enabling new types of interactive workflows. Since our algorithm avoids any global optimization, and its key steps scale linearly with input size, we are able to process extremely large meshes and point clouds, with sizes exceeding several hundred INSERT MODE, 12 characters selected



INSERT MODE, Line 9, Column 1



Daniele Panozzo¹

Olga Sorkine-Hornung¹

³Università dell'Insubria

Introduction

Triangle and quad-dominant meshes are ubiquitously used in computer graphics and CAD applications to represent surfaces, either directly, or as the control grid for higher-order parametric or subdivision surfaces. With the introduction of T-splines [Sederberg et al. 2003] and Dyadic T-Mesh Subdivision [Kovacs et al. 2015], quad-dominant meshes with T-joints (T-meshes) now have similar properties and applications of pure quadrilateral meshes, while being more flexible and naturally supporting the flexible local refinement that is often desired in CAD applications. Meshing surfaces is a challenging problem, and a plethora of methods have been proposed in the past three decades to cope with the increasing quality and scalability requirements of modern applications [Owen 1998; Bommes et al. 2013a]. Semi-regular meshes, which have uniform element



Source	Language	LoC (original)	LoC (IVLA)
Jacobson et al. [2013]	C++	31	8
Sieger and Botsch [2020]	C++	26	9
Alexa [2020]	C++	21	8
Preiner et al. [2019]	Python	15	9
Panozzo et al. [2014]	C++	14	5
Alexa et al. [2020]	C++	12	6
Jakob et al. [2015]	C++	7	4
De Goes and James [2017]	C++	6	1
Jeruzalski et al. [2018]	C++	5	2

Source	Languag
Jacobson et al. [2013]	C++
Sieger and Botsch [2020]	C++
Alexa [2020]	C++
Preiner et al. [2019]	Pythor
Panozzo et al. [2014]	C++
Alexa et al. [2020]	C++
Jakob et al. [2015]	C++
De Goes and James [2017]	C++
Jeruzalski et al. [2018]	C++



Source	Languag
Jacobson et al. [2013]	C++
Sieger and Botsch [2020]	C++
Alexa [2020]	C++
Preiner et al. [2019]	Pythor
Panozzo et al. [2014]	C++
Alexa et al. [2020]	C++
Jakob et al. [2015]	C++
De Goes and James [2017]	C++
Jeruzalski et al. [2018]	C++



• We randomly sampled 100 of all 1987 numbered equations





- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



implementable

- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations





- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



- implementable
- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



implementable

- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



implementable

- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations





- derivations
- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



implementable

- partial derivatives, gradients and unsupported integration
- complicated optimization
- unsupported control flow
- unsupported operators



• We randomly sampled 100 of all 1987 numbered equations



implementable

- complicated optimization
- unsupported control flow
- unsupported operators



• 8 CS PhD students

- 8 CS PhD students
- environment (half C++/Eigen, half Python/NumPy)

Implemented 3 formula using both I\u00f8LA and their preferred programming

- 8 CS PhD students
- environment (half C++/Eigen, half Python/NumPy)

Simple

Given an $n \times n$ matrix A, an n-vector b, and a constant c, compute the quadratic form for an n-vector x:

 $x^T A x$

• Implemented 3 formula using both I \heartsuit LA and their preferred programming

$$(+ b^T x + c)$$

- 8 CS PhD students
- environment (half C++/Eigen, half Python/NumPy)

Medium

Multiply a 3D vertex position v by a weighted average of 4×4 transformation matrices T_i . The corresponding weights are w_i . Assume the vertex position v is already in homogeneous coordinates, which is to say v is a 4-vector.

$$u = \sum_{i}$$

• Implemented 3 formula using both I \heartsuit LA and their preferred programming

 $\sum w_i T_i v$

- 8 CS PhD students
- environment (half C++/Eigen, half Python/NumPy)

Complex

Create an edge-weighted adjacency matrix. Given a set of edges E for a graph of n vertices v_i , create the matrix:

$$A_{ij} = egin{cases} rac{1}{\|v_i - v_j\|} & ext{if } i, j \in E \ 0 & ext{otherwise} \end{cases}$$

• Implemented 3 formula using both I \heartsuit LA and their preferred programming

strongly disagree





Q2: I prefer I LA to the other programming language I used.

neutral	agree	
0%	20%	40%
070	2070	1070

strongly disagree





neutral	agree	
0%	20%	40%

Q3: IULA looks like linear algebra formula I see in papers or on a chalkboard.



User study observations and conclusions

User study observations and conclusions

The average time for each task

sin

I♥LA (minutes) Other (minutes) Significance (*p*)

simple	medium	complex
10	9	12
4	6	12
0.005	0.065	0.862
User study observations and conclusions

The average time for each task

sin

I♥LA (minutes) Other (minutes) Significance (*p*)

Users can accomplish a range of tasks in I\U007LA within 15 minutes

simple	medium	complex
10	9	12
4	6	12
0.005	0.065	0.862

User study observations and conclusions

The average time for each task

sin

- I♥LA (minutes) Other (minutes) Significance (*p*)
- Users can accomplish a range of tasks in IVLA within 15 minutes
- Users perceive that I\U2264 LA looks similar to conventional math

simple	medium	complex
10	9	12
4	6	12
0.005	0.065	0.862

Unsupported operators

 $\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$

Unsupported operators

 $\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$

• Multiple conditions

$$w_{avr}(q) = \frac{1}{\pi} \begin{cases} \frac{\frac{1}{40}(15q^3 - 36q^2 + 40)}{\frac{-3}{4q^3}\left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15}\right) & 1 \le q < 2\\ \frac{3}{4q^3} & q \ge 2. \end{cases}$$



Unsupported operators

$$\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$$

Derivations

$$\sum_{j=1}^{n} P_{i,j} \leq f_0(x_i), \forall i$$
(7)

• Multiple conditions

$$w_{avr}(q) = \frac{1}{\pi} \begin{cases} \frac{\frac{1}{40}(15q^3 - 36q^2 + 40)}{\frac{-3}{4q^3}\left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15}\right) & 1 \le q < 2\\ \frac{3}{4q^3} & q \ge 2. \end{cases}$$



Unsupported operators

$$\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$$

Derivations

$$\sum_{j=1}^{n} P_{i,j} \leq f_0(x_i), \forall i$$
(7)

Multiple conditions

$$w_{\text{avr}}(q) = \frac{1}{\pi} \begin{cases} \frac{1}{40} (15q^3 - 36q^2 + 40) & 0 \le q < 1 \\ \frac{-3}{4q^3} \left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15} \right) & 1 \le q < 2 \\ \frac{3}{4q^3} & q \ge 2 . \end{cases}$$

$$\lambda_{1,2} = 2 \frac{\partial \Psi_5}{\partial I_5},\tag{13}$$





Unsupported operators

$$\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$$

Derivations

$$\sum_{j=1}^{n} P_{i,j} \le f_0(x_i), \, \forall i$$
 (7)

Unsupported optimization

$$\mu(\mathbf{X}) = \max_{\substack{1 \le i \ne j \le n}} \frac{|\mathbf{X}_{.,i}^T \mathbf{X}_{.,j}|}{\|\mathbf{X}_{.,i}\|_2 \|\mathbf{X}_{.,j}\|_2}.$$
(20)

Multiple conditions

$$w_{\text{avr}}(q) = \frac{1}{\pi} \begin{cases} \frac{1}{40} (15q^3 - 36q^2 + 40) & 0 \le q < 1 \\ \frac{-3}{4q^3} \left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15} \right) & 1 \le q < 2 \\ \frac{3}{4q^3} & q \ge 2 . \end{cases}$$

$$\lambda_{1,2} = 2 \frac{\partial \Psi_5}{\partial I_5},\tag{13}$$





Unsupported operators

$$\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$$

Derivations

$$\sum_{j=1}^{n} P_{i,j} \le f_0(x_i), \, \forall i$$
 (7)

Unsupported optimization

$$\mu(\mathbf{X}) = \max_{\substack{1 \le i \ne j \le n}} \frac{|\mathbf{X}_{.,i}^T \mathbf{X}_{.,j}|}{\|\mathbf{X}_{.,i}\|_2 \|\mathbf{X}_{.,j}\|_2}.$$
(20)

Multiple conditions

$$w_{\text{avr}}(q) = \frac{1}{\pi} \begin{cases} \frac{1}{40} (15q^3 - 36q^2 + 40) & 0 \le q < 1 \\ \frac{-3}{4q^3} \left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15} \right) & 1 \le q < 2 \\ \frac{3}{4q^3} & q \ge 2 . \end{cases}$$

Derivatives

$$\lambda_{1,2} = 2 \frac{\partial \Psi_5}{\partial I_5}, \qquad (13)$$

Jers

 $\Psi(x,y,z,t) = \iiint \Phi(k_x,k_y,k_z) e^{2\pi i (k_x x + k_y y + k_z z - ft)} dk_x dk_y dk_z .$ (3)







Unsupported operators

$$\Sigma_i^{\upsilon} = \operatorname{cov}\left(\upsilon_i \cup \mathcal{N}(i)\right) + \sigma_0^2 I,\tag{15}$$

Derivations

$$\sum_{j=1}^{n} P_{i,j} \le f_0(x_i), \, \forall i$$
 (7)

Unsupported optimization

$$\mu(\mathbf{X}) = \max_{\substack{1 \le i \ne j \le n}} \frac{|\mathbf{X}_{.,i}^T \mathbf{X}_{.,j}|}{\|\mathbf{X}_{.,i}\|_2 \|\mathbf{X}_{.,j}\|_2}.$$
(20)

Multiple conditions

$$w_{\text{avr}}(q) = \frac{1}{\pi} \begin{cases} \frac{1}{40} (15q^3 - 36q^2 + 40) & 0 \le q < 1 \\ \frac{-3}{4q^3} \left(\frac{q^6}{6} - \frac{6q^5}{5} + 3q^4 - \frac{8q^3}{3} + \frac{1}{15} \right) & 1 \le q < 2 \\ \frac{3}{4q^3} & q \ge 2 . \end{cases}$$

Derivatives

$$\lambda_{1,2} = 2 \frac{\partial \Psi_5}{\partial I_5}, \qquad (13)$$

Jers

 $\Psi(x,y,z,t) = \iiint \Phi(k_x, k_y, k_z) e^{2\pi i (k_x x + k_y y + k_z z - ft)} dk_x dk_y dk_z .$ (3)







• Support more language features (tensors, automatic differentiation, integration, optimization)

• Support more language features (tensors, automatic differentiation, integration, optimization)

Systematically Differentiating Parametric Discontinuities SAI PRAVEEN BANGARU*, MIT CSAIL JESSE MICHEL*, MIT CSAIL KEVIN MU, MIT CSAIL GILBERT BERNSTEIN, UC Berkeley and MIT CSAIL TZU-MAO LI, MIT CSAIL JONATHAN RAGAN-KELLEY, MIT CSAIL Our Method integrate(x=0 to 1, (x < t) ? 1 oreach x: if x < t: d_out += Our Language (Teg) $[x < t] \mathrm{d}x$ Code Integral with Parametric Discontinuities if x < t out += 1</pre> if x < t d_out + Code Code Traditiona Automated Applications Fig. 1. We propose a language for the automatic differentiation of integrals with discontinuities. Existing auto-diff frameworks require integrals to be discretized into summations prior to differentiation, and therefore lose the derivative contribution from discontinuities. Our method produces a statistically consistent derivative program by introducing integration as a language primitive, which allows us to differentiate discontinuities in continuous space, before discretizing them into summations over discrete samples. Emerging research in computer graphics, inverse problems, and machine language. We introduce integration as a language primitive and account for learning requires us to differentiate and optimize parametric discontinuities. the Dirac delta contribution from differentiating parametric discontinuities These discontinuities appear in object boundaries, occlusion, contact, and in the integrand. We formally define the language semantics and prove the sudden change over time. In many domains, such as rendering and physics correctness and closure under the differentiation, allowing the generation of gradients and higher-order derivatives. We also build a system, TEG, imsimulation, we differentiate the parameters of models that are expressed as integrals over discontinuous functions. Ignoring the discontinuities during plementing these semantics. Our approach is widely applicable to a variety differentiation often has a significant impact on the optimization process. of tasks, including image stylization, fitting shader parameters, trajectory Previous approaches either apply specialized hand-derived solutions, smooth optimization, and optimizing physical designs. out the discontinuities, or rely on incorrect automatic differentiation. We propose a systematic approach to differentiating integrals with dis-CCS Concepts: • Theory of computation \rightarrow Denotational semantics; • continuous integrands, by developing a new differentiable programming Mathematics of computing \rightarrow Differential calculus; Stochastic control and optimization; Probabilistic inference problems; • Computing method-*Both authors contributed equally to this research. $ologies \rightarrow Computer \ graphics; \ Visibility; \ Animation; \ Computer \ vision;$ Authors' addresses: Sai Praveen Bangaru, MIT CSAIL, Cambridge, MA, sbangaru@mit. edu; Jesse Michel, MIT CSAIL, Cambridge, MA, jmmichel@mit.edu; Kevin Mu, MIT CSAIL, Cambridge, MA, kmu@csail.mit.edu; Gilbert Bernstein, UC Berkeley, Berkeley, Derkeley, Berkeley, Marken Marken, Marken, Marken Marken, Marken Marken, Marken Marken, Marken Marken, Ma Modeling and simulation. CA, MIT CSAIL, Cambridge, MA, gilbo@berkelev.edu; Tzu-Mao Li, MIT CSAIL, Cam-Additional Key Words and Phrases: Automatic differentiation, differentiable bridge, MA, tzumao@mit.edu; Jonathan Ragan-Kelley, MIT CSAIL, Cambridge, MA, programming, differentiable graphics, differentiable rendering, differentiable jrk@csail.mit.edu. physics, domain-specific language. Permission to make digital or hard copies of part or all of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation **ACM Reference Format:** For the first page. Copyrights for third-party components of this work must be honored. For all other uses, contact the owner/author(s). 2021 Copyright held by the owner/author(s). Sai Praveen Bangaru, Jesse Michel, Kevin Mu, Gilbert Bernstein, Tzu-Mao Li, and Jonathan Ragan-Kelley. 2021. Systematically Differentiating Parametric Discontinuities. ACM Trans. Graph. 40, 4, Article 107 (August 2021), 17 pages. 0730-0301/2021/8-ART107 https://doi.org/10.1145/3450626.3459775 https://doi.org/10.1145/3450626.3459775 ACM Trans. Graph., Vol. 40, No. 4, Article 107. Publication date: August 2021.

TEG [Bangaru et al. 2021]

Minimization (10%)

Minimization (10%)

$$\Phi(\mathbf{d}) = \min_{\theta} \frac{D_{in}(\theta, \mathbf{d})}{D_{out}(\theta, \mathbf{d})}.$$

$$\hat{\mathbf{x}} = \underset{\mathbf{x}\in\mathcal{M}}{\operatorname{argmin}} \sum_{i} \xi_{i} d(\mathbf{x}, \mathbf{x}_{i})^{2}. \quad (3) \quad \underset{\boldsymbol{\rho}}{\operatorname{minimize}} \left\| \boldsymbol{\tau} - \mathbf{A} \boldsymbol{\rho} \right\|_{2}^{2} + \Gamma\left(\boldsymbol{\rho}\right), \quad \text{s.t. } 0 \leq \boldsymbol{\rho}$$

(10)

$$\max_{\substack{a, \phi_1, \dots, \phi_N \\ \text{s.t. } a^T a = 1}} \sum_{i=1}^N \det\left(p_i, \ \cos\left(\frac{\phi_i}{2}\right) + \sin\left(\frac{\phi_i}{2}\right)a\right).$$
(1)

$$\min_{C} \sum_{i=1}^{N} ||x_i + (R_i - I)C||^2.$$
(3)

$$\mathcal{P}_{opt} = \arg\min_{\mathcal{P}} \Delta \tau_{\mathcal{S}}(\mathcal{P}). \tag{17}$$

$(\hat{\mathbf{I}}, \hat{\mathbf{V}}) = \underset{\mathbf{I}, \mathbf{V}}{\operatorname{arg\,min}} \|\mathbf{J} - \mathbf{\Phi}\mathbf{I}\|_2^2 + R(\mathbf{V}) \quad \text{s.t.} \quad \mathbf{V} = \mathbf{I}.$

$$\min f(\mathbf{v}) = \frac{1}{2} \mathbf{v}^T \mathbf{A} \mathbf{v} - \mathbf{v}^T \mathbf{b}$$
(10)
subject to $\mathbf{J} \mathbf{v} \ge \mathbf{c}_n$.



Minimization (10%)



 $\hat{\mathbf{x}} = \underset{\mathbf{x} \in \mathcal{M}}{\operatorname{argmin}} \sum_{i} \xi_{i} d(\mathbf{x}, \mathbf{x}_{i})^{2} LA:$

 $\max_{\substack{a, \phi_1, \dots, \phi_N \\ \text{s.t. } a^T a = 1}} \sum_{i=1}^{N} \det\left(p_i, \cos\left(\frac{\phi_i}{2}\right) + \sin\left(\frac{\phi}{2}\right) \right) + \sin\left(\frac{\phi}{2}\right) + \sin\left(\frac{\phi}{2}\right) + \exp\left(\frac{\phi}{2}\right) + \exp\left(\frac{\phi}{2}\right$

argmin_(x s.t. $\|x\| > 1$ where $Q \in \mathbb{R}^{3}$

(10)



(1

$$\mathcal{P}_{opt} = \arg\min_{\mathcal{P}} \Delta \tau_{\mathcal{S}}(\mathcal{P}). \qquad (1)$$

$$\mathbf{x} \in \mathbb{R}^{3} \mathbf{)} \mathbf{1/2} \mathbf{x}^{\mathsf{T}} \mathbf{Q} \mathbf{x} + \mathbf{q}^{\mathsf{T}} \mathbf{x} \qquad \text{s.t. } \mathbf{0} \leq \mathbf{\rho}$$

$$\mathbf{x} = \mathbf{I}.$$

$$\min f(\mathbf{v}) = \frac{1}{2} \mathbf{v}^T \mathbf{A} \mathbf{v} - \mathbf{v}^T \mathbf{b}$$
(10)
subject to $\mathbf{J} \mathbf{v} \ge \mathbf{c}_n$.



Integration (8%)

Integration (8%)

$$T_{\rm ir}(\mathbf{x}) = \frac{1}{C} \int_{\mathcal{H}^2} \int_{\mathcal{H}^2} R_{\rm ir}(\mathbf{x}, \boldsymbol{\omega}_i, \boldsymbol{\omega}_o) \,\mathrm{d}\boldsymbol{\omega}_i \,\mathrm{d}\boldsymbol{\omega}_o, \qquad (22)$$

$$\gamma := \int_{-\pi}^{\pi} d(\varphi) \mathbf{c}(\varphi) \, \mathrm{d}\varphi \in \mathbb{C}^{m+1}, \tag{1}$$

$$p(\mathbf{x},k) = -\int \frac{i}{4} H_0^{(2)}(k||\mathbf{x} - \mathbf{y}||)g(\mathbf{y},k) \, d\mathbf{y}.$$
 (14)

$$u_2(x,y) = \frac{e^{ikz}}{i\lambda z} \iint u_1(x',y') e^{\frac{ik}{2z} \{(x-x')^2 + (y-y')^2\}} dx' dy',$$

(4)

$$\widetilde{G} := -\int_{\mathbb{S}^1} \langle \gamma, g \rangle dm \tag{11}$$

$$L_o(\mathbf{x},\boldsymbol{\omega}) = \int_{\partial\Omega} \int_{S^2} L_i(\mathbf{x}',\boldsymbol{\omega}') S(\mathbf{x}',\boldsymbol{\omega}',\mathbf{x},\boldsymbol{\omega}) \,\mathrm{d}\boldsymbol{\omega}' \,\mathrm{d}\mathbf{x}'.$$

$$U(\mathbf{q}) = \int_V \Psi(\mathbf{s}(\mathbf{q})) dV,$$

$$f_{s \to t} = \int F_{\text{Kelvin}} \, \mathrm{d}\boldsymbol{r} = \mu_0 \int \boldsymbol{M}_t(\boldsymbol{r}) \cdot \nabla \boldsymbol{H}_s(\boldsymbol{r}) \, \mathrm{d}\boldsymbol{r}$$
$$= \mu_0 \boldsymbol{m}_t \cdot \int W(\boldsymbol{r} - \boldsymbol{r}_t, h) \nabla H(\boldsymbol{r} - \boldsymbol{r}_s, \boldsymbol{m}_s) \, \mathrm{d}\boldsymbol{r},$$







Integration (8%)

$$T_{\rm ir}(\mathbf{x}) = \frac{1}{C} \int_{\mathcal{H}^2} \int_{\mathcal{H}^2} R_{\rm ir}(\mathbf{x}, \omega_i, \omega_o) \, d\omega_i \, d\omega_o, \qquad (22)$$

$$\widetilde{G} := -\int_{\mathbb{S}^1} \langle \gamma, g \rangle \, dm \qquad (11)$$

$$Y := \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi} \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi \in \mathbb{C}^{m+1} \quad f(\varphi) = \int_{-\pi}^{\pi} d(\varphi) c(\varphi) \, d\varphi = \int_{-\pi}^{\pi} d(\varphi) \, d\varphi = \int_{-\pi}^{\pi} d(\varphi$$

$$f_{s \to t} = \int F_{\text{Kelvin}} \, \mathrm{d}\boldsymbol{r} = \mu_0 \int M_t(\boldsymbol{r}) \cdot \nabla H_s(\boldsymbol{r}) \, \mathrm{d}\boldsymbol{r}$$
$$= \mu_0 \boldsymbol{m}_t \cdot \int W(\boldsymbol{r} - \boldsymbol{r}_t, h) \nabla H(\boldsymbol{r} - \boldsymbol{r}_s, \boldsymbol{m}_s) \, \mathrm{d}\boldsymbol{r},$$







• More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)

- Analyze and support more fields (ML, robotics, CV, physics, scientific computing in general)

• More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)

- Analyze and support more fields (ML, robotics, CV, physics, scientific computing in general)
- Papers of the future

• More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)

- Analyze and support more fields (ML, robotics, CV, physics, scientific computing in general)
- Papers of the future
 - compile an entire paper into a library

More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)

- Analyze and support more fields (ML, robotics, CV, physics, scientific computing in general)
- Papers of the future
 - compile an entire paper into a library
 - improve readability

More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)

- Analyze and support more fields (ML, robotics, CV, physics, scientific computing in general)
- Papers of the future
 - compile an entire paper into a library
 - improve readability

More output languages (PyTorch, TensorFlow, Julia, JavaScript, FORTRAN)



ScholarPhi [Head et al. 2021]

I\U007LA has the potential to greatly benefit the scientific ecosystem

- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas

- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- IVLA can be learned quickly

- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- I\U007LA can be learned quickly
- to readers to implementors

- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- I\U007LA can be learned quickly
- to readers to implementors



- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- I\U007LA can be learned quickly
- to readers to implementors



- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- I\U007LA can be learned quickly
- to readers to implementors





- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- IVLA can be learned quickly
- to readers to implementors





- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- IVLA can be learned quickly
- to readers to implementors




Conclusions

- I\U007LA has the potential to greatly benefit the scientific ecosystem
- IVLA makes it easy to try new ideas
- IVLA can be learned quickly
- to readers to implementors



I\U007LA may reduce translation loss as ideas move from researchers to writers



• Anonymous reviewers for their suggestions

- Anonymous reviewers for their suggestions
- Towaki Takikawa for helpful feedback

ggestions ack

- Anonymous reviewers for their suggestions
- Towaki Takikawa for helpful feedback

Thomas LaToza for a discussion on evaluating programming languages

- Anonymous reviewers for their suggestions
- Towaki Takikawa for helpful feedback

Thomas LaToza for a discussion on evaluating programming languages

- Anonymous reviewers for their suggestions
- Towaki Takikawa for helpful feedback
- Thomas LaToza for a discussion on evaluating programming languages
- Sponsors:
 - Canada Research Chairs Program
 - United States National Science Foundation (IIS-1453018)
 - Adobe







yli69@gmu.edu